

VIRGINIA LAKE L^AT_EX MACROS v4.2

<http://alessio.guglielmi.name/res/vl>

Alessio Guglielmi

University of Bath

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To use this L^AT_EX package: `\usepackage{virginialake}`.

This package is faster and produces better results with `pdflatex`, but can also be used with `latex + dvips`.

If you don't need any graphical elements (most of which have curved lines), you can save some processing time and L^AT_EX resources by doing `\usepackage[noxy]{virginialake}`. In this case, the package does not load X_Y-pic.

Calling the package with `\usepackage[goodsyntax]{virginialake}` produces formulae or structures in the deep-inference style.

Calling the package with `\usepackage[lutzsyntax]{virginialake}` produces formulae or structures in the traditional style, except that it uses square brackets for disjunctions.

To compile this manual from the source make sure that `\write18` is enabled; this can be achieved by invoking `pdflatex` with the `--shell-escape` option.

If T_EX runs out of resources while using Virginia Lake, try invoking `\usepackage{etex}` immediately after `\documentclass`.

For use with Beamer Inside the frame environment in beamer, parentheses and punctuation marks have a special treatment and can cause problems, notably when using square brackets for optional arguments or round brackets for atomic flow coordinates. The solution is to use the option `fragile` for every frame where Virginia Lake macros are used (as in `\begin{frame}[fragile]`). If a frame contains only one slide the option `fragile=singleslide` insures greater speed.

1. SYMBOLS

The following symbols are defined:

1	<code>\one</code>	$c\uparrow$	<code>\cU</code>
0	<code>\zer</code>	$w\downarrow$	<code>\wD</code>
f	<code>\fff</code>	$w\uparrow$	<code>\wU</code>
t	<code>\ttt</code>	$q\downarrow$	<code>\qD</code>
s	<code>\sw</code>	$q\uparrow$	<code>\qU</code>
m	<code>\me</code>	BV	<code>\BV</code>
$i\downarrow$	<code>\iD</code>	SBV	<code>\SBV</code>
$i\uparrow$	<code>\iU</code>	KS	<code>\KS</code>
$c\downarrow$	<code>\cD</code>	SKS	<code>\SKS</code>

2. FORMULAE AND STRUCTURES

Normal formulae:

$$\begin{array}{ll} ((A \vee (B \wedge C) \rightarrow \neg D) \leftrightarrow E) \leftarrow F & ((A \vee \text{lor}(B \vee \text{lan } C) \vee \text{lim} \vee \text{rne } D) \vee \text{ldi } E) \vee \text{lmi } F \\ (A \supset (B \subset C)) & (A \vee \text{ljm } (B \vee \text{lmj } C)) \\ (A \wp (B \otimes (C \triangleleft D))) \rightarrow E & (A \vee \text{lpa } (B \vee \text{lte } (C \vee \text{lse } D))) \vee \text{lli } E \end{array}$$

The command `\vllineartrue` triggers linear logic additive conjunction and disjunction. It is reversed by `\vllinearfalse`, which is the default:

$$((A \oplus (B \& C) \rightarrow \neg D) \leftrightarrow E) \leftarrow F \quad \begin{array}{l} \text{\vllineartrue} \\ ((A \vee (B \wedge C) \rightarrow \neg D) \leftrightarrow E) \leftarrow F \end{array}$$

Logical operators become smaller than usual if `\vlsmallopstrue` is issued. The macro `\vlbin` creates a logical operator: e.g., `\vlse` is `\vlbin\triangleleft`.

Note that the new macros `\lt` and `\gt` are defined, respectively, for `<` and `>` (whose corresponding ASCII characters have special meaning when dealing with formulae).

3. FORMULAE AND STRUCTURES VIA PUNCTUATION AND PARENTHESES

For typesetting formulae, it is possible to use punctuation marks and parentheses instead of macros, as shown below.

The following happens if the package is called without the option `[goodsyntax]`, or if the command `\vlnogoodsyntax` is issued:

$$\begin{array}{ll} A \vee \bar{B} & \{\backslash\text{vls}[A.-B]\} \\ A \vee \bar{B} \vee \dots \vee D \vee E & \{\backslash\text{vls}[A.-B.\backslash\text{vldots}.D.E]\} \\ A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G & \{\backslash\text{vls}[A.(-B.[C;<D;(E;F)>]).G]\} \\ (A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G) & \{\backslash\text{vlsbr}[A.(-B..[C;<D;(E;F)>]).G]\} \\ \{A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G\} & \{\backslash\text{vlscn}[A.(-B.[C;;<D;;(E;;F)>]).G]\} \\ A\{ & A\backslash\text{vlhole} \end{array}$$

The following happens if the package is called with the option `[goodsyntax]`, or if the command `\vlgoodsyntax` is issued:

$$\begin{array}{ll} [A \bar{B}] & \{\backslash\text{vls}[A.-B]\} \\ [A \bar{B} \dots D E] & \{\backslash\text{vls}[A.-B.\backslash\text{vldots}.D.E]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\backslash\text{vls}[A.(-B.[C;<D;(E;F)>]).G]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\backslash\text{vlsbr}[A.(-B..[C;<D;(E;F)>]).G]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\backslash\text{vlscn}[A.(-B.[C;;<D;;(E;;F)>]).G]\} \\ A\{ & A\backslash\text{vlhole} \end{array}$$

The following happens if the package is called with the option `[lutzsyntax]`, or if the command `\vllutzsyntax` is issued:

$$\begin{array}{ll} A \vee \bar{B} & \{\backslash\text{vls}[A.-B]\} \\ A \vee \bar{B} \vee \dots \vee D \vee E & \{\backslash\text{vls}[A.-B.\backslash\text{vldots}.D.E]\} \\ A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G & \{\backslash\text{vls}[A.(-B.[C;<D;(E;F)>]).G]\} \\ [A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G] & \{\backslash\text{vlsbr}[A.(-B..[C;<D;(E;F)>]).G]\} \\ \{A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G\} & \{\backslash\text{vlscn}[A.(-B.[C;;<D;;(E;;F)>]).G]\} \\ A\{ & A\backslash\text{vlhole} \end{array}$$

The command `\vlsmallbrackets` corrects a possible problem with the size of brackets:

$$\begin{array}{l} (a \vee b) \wedge (\bar{a} \vee \bar{b}) \quad \{\backslash\text{vls}([a.b].[-a.-b])\} \\ \hline (a \vee b) \wedge (\bar{a} \vee \bar{b}) \quad \{\backslash\text{vls}\smallbrackets \\ \quad \backslash\text{vls}([a.b].[-a.-b])\} \end{array}$$

This also holds for the alternative syntax style:

$$\begin{array}{l} ([a b] [\bar{a} \bar{b}]) \quad \{\backslash\text{vls}([a.b].[-a.-b])\} \\ \hline ([a b] [\bar{a} \bar{b}]) \quad \{\backslash\text{vls}\smallbrackets \\ \quad \backslash\text{vls}([a.b].[-a.-b])\} \end{array}$$

The command `\vlnosmallbrackets` undoes `\vlsmallbrackets`.

The `\vls` macro works by redefining ‘, ‘[’, ‘]’, ‘(’ and ‘)’. This, of course, can cause several problems. In practice, they are rare, except when `\vls` is an argument of a macro. In this case, one needs to ‘update’ the macro by calling `\vlupdate`. Compare

```
[a.b] \newcommand{\vlttest}[1]{#1}
      \vlttest{\vls[a.b]}
```

with

```
a \ve b \newcommand{\vlttest}[1]{#1}
      \vlupdate\vlttest
      \vlttest{\vls[a.b]}
```

This mechanism, in principle, works with every macro.

Sometimes it is not possible to use `\vlupdate`. For example, in the following situation (requiring the `amsmath` package):

$\beta_k = \bigwedge_{k \leq i \leq n} [-c_i . -d_i] \quad \text{for } 1 < k \leq n,$ $\gamma_i = (\beta_{i+1} . c_i) \quad \text{for } 1 \leq i < n,$ $\delta_i = (\beta_{i+1} . d_i) \quad \text{for } 1 \leq i < n,$ $\gamma_n = c_n,$ $\delta_n = d_n.$	<pre>\begin{align*} \beta_k & \&= \bigwedge_{k \leq i \leq n} & \vls[-c_i . -d_i] & & & \text{\textfor } \\$1 < k \le n \\$,} \\ \gamma_i & = \vls(\beta_{i+1} . c_i) & & \text{\textfor } \\$1 \le i < n \\$,} \\ \delta_i & = \vls(\beta_{i+1} . d_i) & & \text{\textfor } \\$1 \le i < n \\$,} \\ \gamma_n & = c_n, & & \\ \delta_n & = d_n. & & \\ \end{align*}</pre>
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The solution is to use the macros `\vlstore`, `\vlread` and `\lt` as follows:

$\beta_k = \bigwedge_{k \leq i \leq n} \bar{c}_i \ve \bar{d}_i \quad \text{for } 1 < k \leq n,$ $\gamma_i = \beta_{i+1} \wedge c_i \quad \text{for } 1 \leq i < n,$ $\delta_i = \beta_{i+1} \wedge d_i \quad \text{for } 1 \leq i < n,$ $\gamma_n = c_n,$ $\delta_n = d_n.$	<pre>\vlstore{% \beta_k &= \bigwedge_{k \leq i \leq n} & \vls[-c_i . -d_i] & & & \text{\textfor } \\$1 \lt k \le n \\$,} \\ \gamma_i &= \vls(\beta_{i+1} . c_i) & & \text{\textfor } \\$1 \le i \lt n \\$,} \\ \delta_i &= \vls(\beta_{i+1} . d_i) & & \text{\textfor } \\$1 \le i \lt n \\$,} \\ \gamma_n &= c_n, & & \\ \delta_n &= d_n \vldot & & \\ } \begin{align*} \vlread \end{align*}</pre>
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Note the use of `\vldot` in the place of `.` in the last line of the display. The commands `\vldot` for `.` and `\vlsqbrl` for `[` and `\vlsqbrl` for `]` and `\vlrobrl` for `(` and `\vlrobrl` for `)` are provided.

There’s the command `\vlnos`, which sets ‘, ‘[’, ‘]’, ‘(’ and ‘)’ to their normal behaviour, for example:

```
a \wedge ([b \wedge c]) \vls(a. {\vlnos (. {\vls(b.c)} .)})
```

4. DEEP INFERENCE AND SEQUENT CALCULUS DERIVATIONS

The following macros, by default, accept the commands for formulae and structures, as shown before. However, this might create problems in some circumstances, because the characters ‘, ‘;’, ‘[’, ‘]’, ‘(’, ‘)’, ‘<’, ‘>’ and ‘-’, are defined in a special way. To avoid such problems, in case the commands for formulae/structures are not wanted, you can use the command `\vlnostruressyntax`: every successive invocation of the commands for derivations does not redefine any character.

$\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vlin{\rho}\{\star\}\{B\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vlid{\rho}\{\star\}\{B\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vliq{\rho}\{\star\}\{B\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vlinf{\rho}\{\star\}\{B\}\{A\}}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vlidf{\rho}\{\star\}\{B\}\{A\}}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vliqf{\rho}\{\star\}\{B\}\{A\}}$	$\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vliin{\rho}\{\star\}\{C\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vliid{\rho}\{\star\}\{C\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \begin{array}{l} \text{\vlderivation} \\ \text{\vliiq{\rho}\{\star\}\{C\}\{A\}} \\ \text{\vlhy} \end{array}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vliinf{\rho}\{\star\}\{C\}\{A\}\{B\}}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vliidf{\rho}\{\star\}\{C\}\{A\}\{B\}}$ <hr/> $\frac{\rho \frac{A}{B} \star}{\rho \frac{A}{B} \star} \text{\vliiqf{\rho}\{\star\}\{C\}\{A\}\{B\}}$
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$\rho \frac{A \ B \ C}{D} \star$	\backslash vlderivation { \backslash vliiin{\rho}{\star}{D}{ \backslash vlhy {A} \backslash vlhy { \backslash vlhy {B} \backslash vlhy { \backslash vlhy {C}}
$\rho \frac{A \ B \ C}{D} \star$	\backslash vlderivation { \backslash vliiid{\rho}{\star}{D}{ \backslash vlhy {A} \backslash vlhy { \backslash vlhy {B} \backslash vlhy { \backslash vlhy {C}}
$\rho \frac{A \ B \ C}{D} \star$	\backslash vlderivation { \backslash vliiiq{\rho}{\star}{D}{ \backslash vlhy {A} \backslash vlhy { \backslash vlhy {B} \backslash vlhy { \backslash vlhy {C}}
$\rho \frac{A \ B \ C}{D} \star$	\backslash vliiin{\rho}{\star}{D}{A}{B}{C}
$\rho \frac{A \ B \ C}{D} \star$	\backslash vliiid{\rho}{\star}{D}{A}{B}{C}
$\rho \frac{A \ B \ C}{D} \star$	\backslash vliiiq{\rho}{\star}{D}{A}{B}{C}

$\rho \frac{A \ B \ C \ D}{E} \star$	\vlderivation { $\text{\vliiiin{\rho}\{\star\}\{E\}\{$ \vlhy {A} \vlhy {B} \vlhy {C} \vlhy {D}}
$\rho \frac{A \ B \ C \ D}{\dots\dots\dots E} \star$	\vlderivation { $\text{\vliiiid{\rho}\{\star\}\{E\}\{$ \vlhy {A} \vlhy {B} \vlhy {C} \vlhy {D}}
$\rho \frac{A \ B \ C \ D}{\overline{\overline{E}}} \star$	\vlderivation { $\text{\vliiiiq{\rho}\{\star\}\{E\}\{$ \vlhy {A} \vlhy {B} \vlhy {C} \vlhy {D}}
$\rho \frac{A \ B \ C \ D}{E} \star$	$\text{\vliiiinf{\rho}\{\star\}\{E\}\{A\}\{B\}\{C\}\{D\}}$
$\rho \frac{A \ B \ C \ D}{\dots\dots\dots E} \star$	$\text{\vliiiidf{\rho}\{\star\}\{E\}\{A\}\{B\}\{C\}\{D\}}$
$\rho \frac{A \ B \ C \ D}{\overline{\overline{E}}} \star$	$\text{\vliiiiqf{\rho}\{\star\}\{E\}\{A\}\{B\}\{C\}\{D\}}$

$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl derivation { \vl iiii in{\rho}{\star}{F}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} </pre>
<hr/>	
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl derivation { \vl iiii id{\rho}{\star}{F}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} </pre>
<hr/>	
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl derivation { \vl iiii iq{\rho}{\star}{F}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} </pre>
<hr/>	
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl iiii inf{\rho}{\star}{F}{A}{B}{C}{D}{E} </pre>
<hr/>	
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl iiii idf{\rho}{\star}{F}{A}{B}{C}{D}{E} </pre>
<hr/>	
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre> \vl iiii iqf{\rho}{\star}{F}{A}{B}{C}{D}{E} </pre>

$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl derivation { \vl iiii in{\rho}{\star}{G}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} { \vl hy {F}}}} </pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl derivation { \vl iiii id{\rho}{\star}{G}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} { \vl hy {F}}}} </pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl derivation { \vl iiii iq{\rho}{\star}{G}{ \vl hy {A}} { \vl hy {B}} { \vl hy {C}} { \vl hy {D}} { \vl hy {E}} { \vl hy {F}}}} </pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl iiii inf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F} </pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl iiii idf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F} </pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre> \vl iiii iqf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F} </pre>

The commands `\vlsmallleftlabels` and `\vlnosmallleftlabels` control the size of labels at the left of — (the ρ s above); the default is small size.

$\Pi \parallel S$	<code>\vlderivation</code>	{
$B \vee C$	<code>\vlpr{\Pi}{\cal S}{\vls[B.C]}</code>	}
$\Pi \parallel S$	<code>\vlderivation</code>	{
$B \vee C$	<code>\vlpd{\Pi}{\cal S}{\vls[B.C]}</code>	}
$\Pi \parallel S$	<code>\vlderivation</code>	{
$B \vee C$	<code>\vlpf{\Pi}{\cal S}{\vls[B.C]}{\the\toks0}</code>	}
$\Pi \parallel S$	<code>\vlproof{\Pi}{\cal S}{\vls[B.C]}</code>	
$\Pi \parallel S$	<code>\vlproofd{\Pi}{\cal S}{\vls[B.C]}</code>	
A	<code>\vlderivation</code>	{
$\Delta \parallel S$	<code>\vlde{\Delta}{\cal S}{\vls[B.C]}{</code>	
$B \vee C$	<code>\vlhy {A }}}</code>	
A	<code>\vlderivation</code>	{
$\Delta \parallel S$	<code>\vldd{\Delta}{\cal S}{\vls[B.C]}{</code>	
$B \vee C$	<code>\vlhy {A }}}</code>	
A	<code>\vlderivation</code>	{
$\Delta \parallel S$	<code>\vldf{\Delta}{\cal S}{\vls[B.C]}{</code>	
$B \vee C$	<code>\vlhy {A }{\the\toks0}</code>	}
A	<code>\vllder{\Delta}{\cal S}{\vls[B.C]}{A}</code>	
A	<code>\vllderd{\Delta}{\cal S}{\vls[B.C]}{A}</code>	

$\Pi \parallel_S$ $\rho \frac{A}{B \vee C}$ note	<code>\vlderivation</code> <code>\vlin{\rho}{\rm note}{\vls[B.C]}{</code> <code>\vlpr{\Pi }{\cal S }{A }</code> <code>}</code>
A $\Delta \parallel_S$ $\rho \frac{B}{C}$	<code>\vlderivation</code> <code>\vlin{\rho }{ }</code> <code>\vlde{\Delta}{\cal S}{B}{</code> <code>\vlhy</code> <code>{A}}}</code>
A $\Delta \parallel_{S''}$ B $\Delta' \parallel_{S'}$ C $\Delta'' \parallel_S$ D	<code>\vlderivation</code> <code>\vlde{\Delta''}{\cal S }{D} {</code> <code>\vlde{\Delta' }{\cal S' }{C} {</code> <code>\vlde{\Delta }{\cal S''}{B}{</code> <code>\vlhy</code> <code>{A}}}}</code>
$\frac{A}{\rho B}$ $\Delta \parallel$ C $\Delta' \parallel$ D	<code>\vlderivation</code> <code>\vlde{\Delta'}{ }{D} {</code> <code>\vlde{\Delta }{ }{C} {</code> <code>\vlin{\rho }{ }{B}{</code> <code>\vlhy</code> <code>{A}}}}</code>
A $\rho \overline{\overline{B}}$ $\rho' \overline{\overline{C}}$ $\rho'' \overline{\overline{D}}$	<code>\vlderivation</code> <code>\vlin{\rho''}{\quad.}{D} {</code> <code>\vliq{\rho' }{ }{C} {</code> <code>\vlid{\rho }{ }{B}{</code> <code>\vlhy</code> <code>{A}}}}</code>
$\frac{A \wedge A'}{\rho \overline{\overline{B \wedge B'}}$ $\Delta \parallel$ $\frac{C \wedge C'}{\rho' \overline{\overline{C \wedge C'}}$ $\Delta' \parallel$ $\frac{D \wedge D'}{\rho'' \overline{\overline{D \wedge D'}}$ *	<code>\vlderivation</code> <code>\vliiq{\rho }{\star}{\vls(E.E')}</code> <code>\vlde {\Delta'}{ }{\vls(D.D')}</code> <code>\vlde {\Delta }{ }{\vls(C.C')}</code> <code>\vlin {\rho }{ }{\vls(B.B')}</code> <code>\vlhy</code> <code>{\vls(A.A')}}}}</code> <code>{</code> <code>\vlin {\rho'' }{\star}{\vls(D.D')}</code> <code>\vlin {\rho' }{ }{\vls(C.C')}</code> <code>\vlin {\rho }{ }{\vls(B.B')}</code> <code>\vlhy</code> <code>{\vls(A.A')}}}}</code> <code>}</code>
$\frac{E \wedge E'}{\rho \overline{\overline{E \wedge E'}}$ *	<code>\vlderivation</code> <code>\vlin {\rho'' }{\star}{\vls(D.D')}</code> <code>\vlin {\rho' }{ }{\vls(C.C')}</code> <code>\vlin {\rho }{ }{\vls(B.B')}</code> <code>\vlhy</code> <code>{\vls(A.A')}}}}</code> <code>}</code>

Note `\vldots` in

γ \parallel α_1 $\tau \frac{\alpha_1 \sigma_1}{\alpha_1 \sigma_1}$ \parallel \vdots \parallel α_n $\rho \frac{\alpha_n \sigma_n}{\alpha_n \sigma_n}$ \parallel β	<code>\vlderivation</code> <code>\vlde{ }{\beta }</code> <code>\vlin{\rho}{\alpha_n \sigma_n}</code> <code>\vlde{ }{\alpha_n }</code> <code>\vlde{ }{\vldots }</code> <code>\vlin{\tau}{\alpha_1 \sigma_1}</code> <code>\vlde{ }{\alpha_1 }</code> <code>\vlhy</code> <code>{\gamma }</code> <code>}</code>
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The command `\vlx` provides some extra space, as in

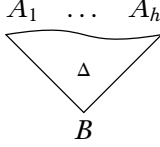
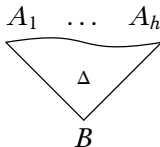
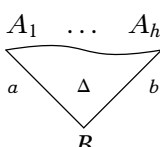
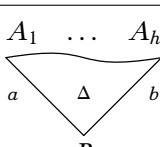
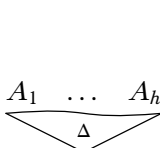

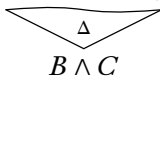
$$\xi \left\{ \begin{array}{c} \Psi \parallel \\ \frac{t}{a \vee \bar{a}} \wedge (\bar{a} \vee \bar{a}) \\ \frac{a}{a \wedge a} \wedge (\bar{a} \vee \bar{a}) \\ \frac{a \wedge ((a \wedge \bar{a}) \vee \bar{a})}{a \wedge \bar{a}} \vee \bar{a} \\ \frac{a \wedge \bar{a}}{f} \vee \frac{a \wedge \bar{a}}{f} \\ \Psi \parallel \\ \beta \end{array} \right\}$$

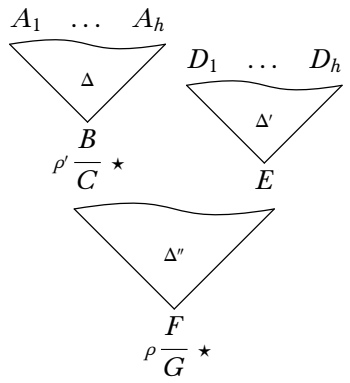
```

\vlx \vlx
{-a}
]
{\vls(\vlinf{}){\vls[a.-a]}t. [-a.-a]}
%-----
\right\}
\vlhy \alpha
}}

```

The following use Xy-pic (see at the beginning of this manual for instructions).

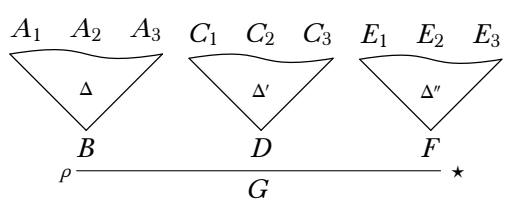
	<pre> \lderivation { \ltr{\Delta}{B }{ \lhy {A_1 }} { \lhy {\dots}} { \lhy {A_h }}}</pre>
	<pre> \ltreeder{\Delta}{B}{A_1}{\dots}{A_h}</pre>
	<pre> \lderivation { \ltrl{\Delta}{a}{b}{B }{ \lhy {A_1 }} { \lhy {\dots}} { \lhy {A_h }}}</pre>
	<pre> \ltreederl{\Delta}{a}{b}{B}{A_1}{\dots}{A_h}</pre>
	<pre> {\vlnostruressyntax \lderivation { \ltrf{\Delta}{B }{ \lhy {A_1 }} { \lhy {\dots}} { \lhy {A_h }} {0.5}}</pre>
	<pre> \toks0={0.5} \lderivation { \ltrf{\Delta}{\vls(B.C)}{ \lhy {A_1 }} { \lhy {\dots }} { \lhy {A_h }} {\the\toks0}}</pre>
	<pre> \toks0={2.5} \lderivation { \ltrlf{\Delta}{a}{b}{\vls(B.C)}{ \lhy {A_1 }} { \lhy {\dots }} { \lhy {A_h }} {\the\toks0}}</pre>



```

\dimen0=-1pc
\vlderivation
\vltr{\rho }{\star}{G } {
\vltr{\Delta''} {F } {
\vltr{\rho' }{\star}{C } {
\vltr{\Delta } {B } {
\vlhy {A_1 } {
\vlhy {\dots } {
\vlhy {A_h } {}}}
\vlhy {\kern\dimen0} {
\vltr{\Delta' } {E } {
\vlhy {D_1 } {
\vlhy {\dots } {
\vlhy {D_h } {}}}

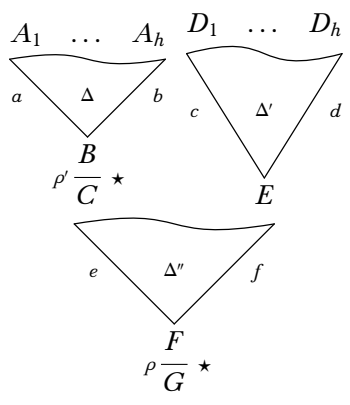
```



```

\vlderivation {
\vltr{\rho }{\star}{G } {
\vltr {\Delta } {B }{
\vlhy {A_1}}
{
\vlhy {A_2}}
{
\vlhy {A_3}}}
\vltr {\Delta' } {D }{
\vlhy {C_1}}
{
\vlhy {C_2}}
{
\vlhy {C_3}}}
\vltr {\Delta''} {F }{
\vlhy {E_1}}
{
\vlhy {E_2}}
{
\vlhy {E_3}}}}

```



```

{\vlnostruressyntax
\vlderivation
\vlrh {rho }{\star}{G } {
\vltrl {\Delta'}{e}{f} {F } {
\vlrh {\rho'} {\star}{C } {
\vltrl {\Delta }{a}{b} {B }{
\vlhy {A_1 } }
{
\vlhy {\dots } }
{
\vlhy {A_h } }}}}
{
\vlhy {\kern-1pc} }
{
\vltrlf{\Delta' }{c}{d} {E }{
\vlhy {D_1 } }
{
\vlhy {\dots } }
{
\vlhy {D_h } }
{1.6}}}}

```

5. OPEN DEDUCTION

The following macros for open deduction do not reverse the order of formulae.



















$\rho \frac{A}{B \vee C} \star$	$\{A\}$	$\backslash\text{od}\{\backslash\text{odi}\{\backslash\text{odh}\}$
$\rho \frac{A}{B \vee C} \star$	$\{A\}$	$\backslash\text{od}\{\backslash\text{odo}\{\backslash\text{odh}\}$
$\rho \frac{A}{B \vee C} \star$	$\{A\}$	$\backslash\text{od}\{\backslash\text{odI}\{\backslash\text{odh}\}$
$\rho \frac{A}{B \vee C} \star$	$\backslash\text{odn}\{A\}\{\rho\}\{[B.C]}\{\backslash\text{star}\}$	
$\rho \frac{A}{B \vee C} \star$	$\backslash\text{odt}\{A\}\{\rho\}\{[B.C]}\{\backslash\text{star}\}$	
$\rho \frac{A}{B \vee C} \star$	$\backslash\text{odN}\{A\}\{\rho\}\{[B.C]}\{\backslash\text{star}\}$	
$\Pi \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{odp}\}$
$\Pi \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{odP}\}$
$\Pi \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{odpx}\}$
$\Pi \parallel s$	$\backslash\text{odr}\{\Pi\}\{[B.C]}\{\backslash\text{cal S}\}$	
$\Pi \parallel s$	$\backslash\text{odR}\{\Pi\}\{[B.C]}\{\backslash\text{cal S}\}$	
$\Delta \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{odd}\{\backslash\text{odh}\}$
$\Delta \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{odD}\{\backslash\text{odh}\}$
$\Delta \parallel s$	$B \vee C$	$\backslash\text{od}\{\backslash\text{oddX}\{\backslash\text{odh}\}$
$\Delta \parallel s$	$\backslash\text{odv}\{A\}\{\Delta\}\{[B.C]}\{\backslash\text{cal S}\}$	
$\Delta \parallel s$	$\backslash\text{odV}\{A\}\{\Delta\}\{[B.C]}\{\backslash\text{cal S}\}$	

$\Pi \parallel S$ $\rho \frac{A}{B \vee C}$ note	<pre>\odf\odi{\odp {\Pi }{A }{\cal S } } {\rho}{[B.C]}{\rm note}}</pre>
$\Delta \parallel S$ $\rho \frac{A}{B}$ $\rho \frac{B}{C}$	<pre>\odf\odi{\odd{\odh {A} } } {\Delta}{B}{\cal S } } {\rho }{C}{ }</pre>
$\Delta \parallel S''$ $\Delta' \parallel S'$ $\Delta'' \parallel S$ D	<pre>\odf\odd{\odd{\odd{\odh {A} } } } {\Delta }{B}{\cal S ''} } {\Delta' }{C}{\cal S ' } } {\Delta''}{D}{\cal S }}</pre>
$\rho \frac{A}{B}$ $\Delta \parallel C$ $\Delta' \parallel D$	<pre>\odf\odd{\odd{\odi{\odh {A} } } } {\rho }{B}{ } } {\Delta }{C}{ } } {\Delta' }{D}{ }</pre>
$\rho \frac{A}{B}$ $\rho \frac{B}{C}$ $\rho \frac{C}{D}$	<pre>\odf\odi{\odI{\odof{\odh {A} } } } {\rho }{B}{ } } {\rho' }{C}{ } } {\rho''}{D}{\quad.}}</pre>
$A \wp B$ $\psi_1 \parallel$ $a \wp b \wp \left[\begin{array}{c} \bar{b} \otimes b \\ \perp \end{array} \right] \wp (\bar{a} \otimes \bar{b})$	<pre>\odV{[A;B]} {\psi_1}{\odbox{[a;b ; \odn{(-b;b)} \iU\bot {} ; (-a;-b)] }}}</pre>
$\frac{a}{a \wedge t} \vee \frac{\bar{a}}{t \wedge \bar{a}}$ $(a \vee t) \wedge \frac{t \vee \bar{a}}{\bar{a} \vee t}$ $\frac{(a \vee t) \wedge \bar{a}}{\bar{a} \wedge (a \vee t)} \vee t$ $\frac{\bar{a} \wedge a}{f} \vee t$	<pre>\odbackgroundtrue {\ttt \iD{[\odn{ a } ={([a.\ttt])}{]}\odn{ -a } ={(\ttt.-a)}{]} } \me{([a.\ttt]..\odn{[\ttt.-a]} =[[-a.\ttt]}{]) }{]} \sw{[\odf\odi{\odi{\odh {([a.\ttt].-a) } } =[(-a.[a.\ttt]) }{]} \sw{[\odn{(-a.a)} \iU{\fff }{]}\ttt]}{]}\ttt]}{]} }</pre>




You can suppress or enable the background colour and the frames, respectively, with the commands `\odbackgroundfalse` or `\odbackgroundtrue`, and `\odframefalse` or `\odframetrue`.

6. ATOMIC FLOWS

The following uses \Xypic (see at the beginning of this manual for instructions).

	<code>\afid abcd</code>		<code>\afiu abcd</code>
	<code>\afidc abcd{Red}{Green}</code>		<code>\afiuc abcd{Red}{Green}</code>
	<code>\afidx abcd45</code>		<code>\afiux abcd32</code>
	<code>\afidxc abcd45{Red}{Green}</code>		<code>\afiuxc abcd32{Red}{Green}</code>
	<code>\afidn</code>		<code>\afiun</code>
	<code>\afId abcd</code>		<code>\afIu abcd</code>
	<code>\afIdc abcd{Red}{Green}</code>		<code>\afIuc abcd{Red}{Green}</code>
	<code>\afIdx abcd45</code>		<code>\afIux abcd32</code>
	<code>\afIdxc abcd45{Red}{Green}</code>		<code>\afIuxc abcd32{Red}{Green}</code>

The three macros `\aftrim`, `\aftrimabove` and `\aftrimbelow` trim extra space above and below diagrams including `\afiu` and similar macros.

	<code>\aftrim{\af{(0,0)*{\afwun}}}</code>
	<code>\aftrimabove{\af{(0,0)*{\afid{}{}{}}}}</code>
	<code>\aftrimbelow{\af{(0,0)*{\afwu{}{}}}}</code>

Use the macros `\afraise` and `\aflower`, typically with one atomic flow as an argument, to raise or lower a bit the atomic flow (see examples in the following).

∇ \afwd {}{}

∇ \afwd ab
 $a|b$

∇ \afwdc ab{Red}
 $a|b$

∇ \afwdn

∇ \afWd {}{}

∇ \afWd ab
 $a||b$

∇ \afWdc ab{Red}
 $a||b$

\blacktriangledown \afwds {}{}

\blacktriangledown \afwds ab
 $a|b$

\blacktriangledown \afwdsc ab{Red}
 $a|b$

\blacktriangledown \afwdsn

\blacktriangledown \afWds {}{}

\blacktriangledown \afWds ab
 $a||b$

\blacktriangledown \afWdsc ab{Red}
 $a||b$

\blacktriangle \afwu {}{}

\blacktriangle \afwu ab
 $a|b$

\blacktriangle \afwuc ab{Red}
 $a|b$

\blacktriangle \afwun

\blacktriangle \afWu {}{}

\blacktriangle \afWu ab
 $a||b$

\blacktriangle \afWuc ab{Red}
 $a||b$

\blacktriangleup \afwus {}{}

\blacktriangleup \afwus ab
 $a|b$

\blacktriangleup \afwusc ab{Red}
 $a|b$

\blacktriangleup \afwusn

\blacktriangleup \afWus {}{}

\blacktriangleup \afWus ab
 $a||b$

\blacktriangleup \afWusc ab{Red}
 $a||b$

\bullet \afdd {}{}

 \bullet \afdd ab
 $a|b$

 \bullet \afddc ab{Red}
 $a|b$

 \bullet \afDd {}{}
 $a||b$

 \bullet \afDd ab
 $a||b$

 \bullet \afDdc ab{Red}
 $a||b$

\bullet \afdu {}{}

 $a|b$
 \bullet \afdu ab

 $a|b$
 \bullet \afduc ab{Red}

 \bullet \afDu {}{}

 $a||b$
 \bullet \afDu ab

 $a||b$
 \bullet \afDuc ab{Red}

\bullet \afd

↓ \afad {}{}

$a|b$
↓ \afad ab

$a|b$
↓ \afadc ab{Red}

▽ \afadn

⇓ \afAd {}{}

$a||b$
⇓ \afAd ab

$a||b$
⇓ \afAdc ab{Red}

↓ \afads {}{}

$a|b$
↓ \afads ab

$a|b$
↓ \afadsc ab{Red}

▽ \afadsn

⇓ \afAds {}{}

$a||b$
⇓ \afAds ab

$a||b$
⇓ \afAdsc ab{Red}

↑ \afau {}{}

$a|b$
↑ \afau ab

$a|b$
↑ \afauc ab{Red}

△ \afaun

⇑ \afAu {}{}

$a||b$
⇑ \afAu ab

$a||b$
⇑ \afAuc ab{Red}

↑ \afaus {}{}

$a|b$
↑ \afaus ab


$a|b$
↑ \afausc ab{Red}


▲ \afausn


⇑ \afAus {}{}


$a||b$
⇑ \afAus ab


$a||b$
⇑ \afAusc ab{Red}


 \afcd {}{}{}{}{}{}


 \afcd abcdef


 \afcdc abcdef{Red}{Green}{Blue}


 \afcdn abcd


 \afcdnc abcd{Red}{Green}


 \afCd {}{}{}{}{}{}


 \afCd abcdef


 \afCdc abcdef{Red}{Green}{Blue}


 \afCdn abcd


 \afCdnc abcd{Red}{Green}


 \afcds {}{}{}{}{}{}


 \afcds abcdef


 \afcdsc abcdef{Red}{Green}{Blue}

 \afcdsn abcd


 \afcdsnc abcd{Red}{Green}


 \afCds {}{}{}{}{}{}


 \afCds abcdef


 \afCdsc abcdef{Red}{Green}{Blue}


 \afCdsn abcd


 \afCdsnc abcd{Red}{Green}


 \afcu {}{}{}{}{}{}


 \afcu abcdef


 \afcuc abcdef{Red}{Green}{Blue}


 \afcun abcd


 \afcunc abcd{Red}{Green}


 \afCu {}{}{}{}{}{}


 \afCu abcdef


 \afCuc abcdef{Red}{Green}{Blue}


 \afCun abcd


 \afCunc abcd{Red}{Green}


 \afcus {}{}{}{}{}{}


 \afcus abcdef


 \afcusc abcdef{Red}{Green}{Blue}


 \afcun abcd


 \afcusnc abcd{Red}{Green}

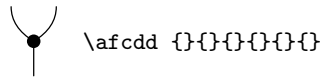
 \afCus {}{}{}{}{}{}

 \afCus abcdef

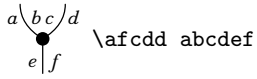
 \afCusc abcdef{Red}{Green}{Blue}

 \afCun abcd

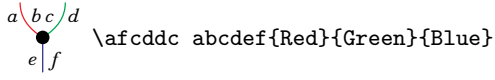
 \afCusnc abcd{Red}{Green}



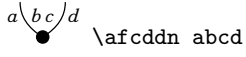
`\afcdd {}{}{}{}{}{}`



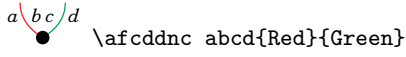
`\afcdd abcdef`



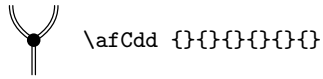
`\afcddc abcdef{Red}{Green}{Blue}`



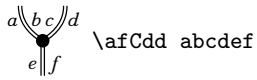
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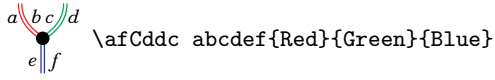
`\afcddnc abcd{Red}{Green}`



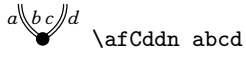
`\afCdd {}{}{}{}{}{}`



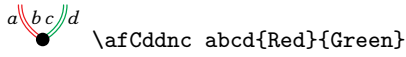
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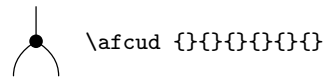
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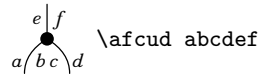
`\afCddn abcd`



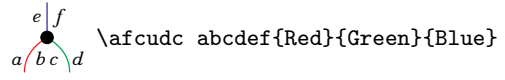
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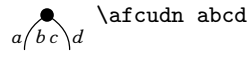
`\afcud {}{}{}{}{}{}`



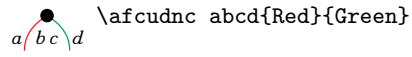
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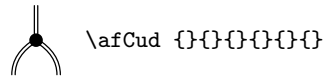
`\afcudc abcdef{Red}{Green}{Blue}`



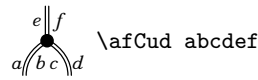
`\afcudn abcd`



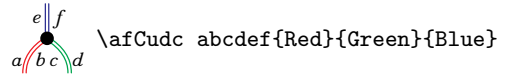
`\afcudnc abcd{Red}{Green}`



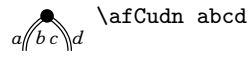
`\afCud {}{}{}{}{}{}`



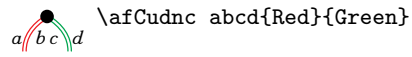
`\afCud abcdef`



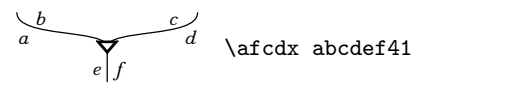
`\afCudc abcdef{Red}{Green}{Blue}`



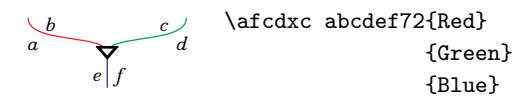
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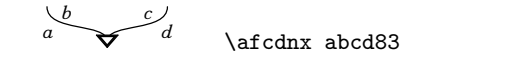
`\afCudnc abcd{Red}{Green}`



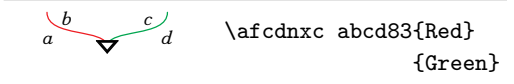
`\afcdx abcdef41`



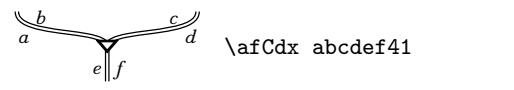
`\afcdxc abcdef72{Red}`
`{Green}`
`{Blue}`



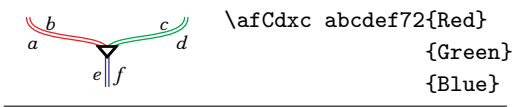
`\afcdnx abcd83`



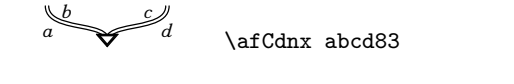
`\afcdnxc abcd83{Red}`
`{Green}`



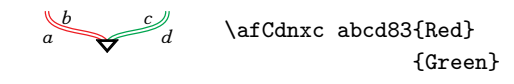
`\afCdx abcdef41`



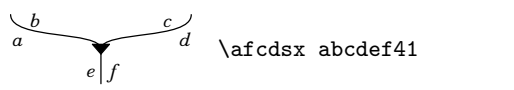
`\afCdx abcdef72{Red}`
`{Green}`
`{Blue}`



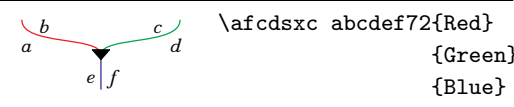
`\afCdnx abcd83`



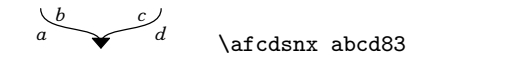
`\afCdnxc abcd83{Red}`
`{Green}`



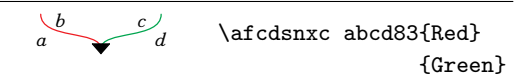
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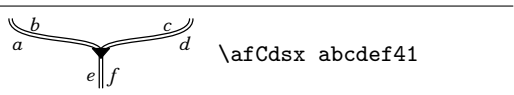
`\afcdsxc abcdef72{Red}`
`{Green}`
`{Blue}`



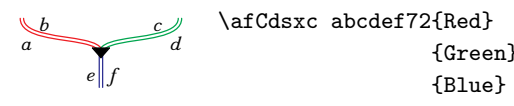
`\afcdsnx abcd83`



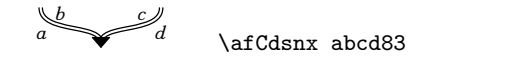
`\afcdsnxc abcd83{Red}`
`{Green}`



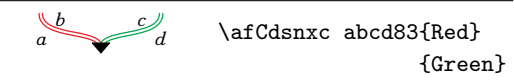
`\afCdsx abcdef41`



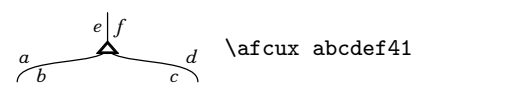
`\afCdsxc abcdef72{Red}`
`{Green}`
`{Blue}`



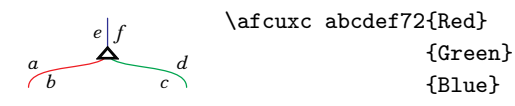
`\afCdsnx abcd83`



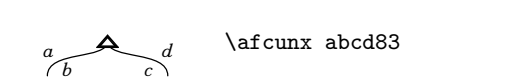
`\afCdsnxc abcd83{Red}`
`{Green}`



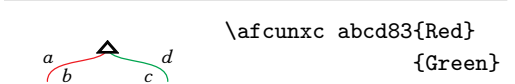
`\afcux abcdef41`



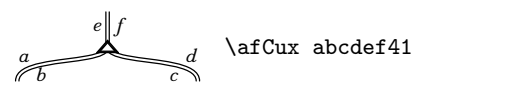
`\afcux abcdef72{Red}`
`{Green}`
`{Blue}`



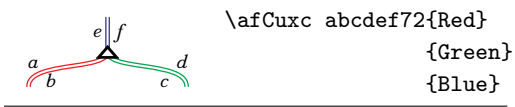
`\afcunx abcd83`



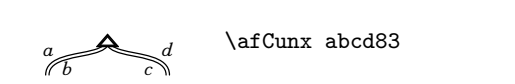
`\afcunxc abcd83{Red}`
`{Green}`



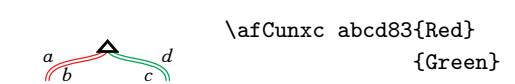
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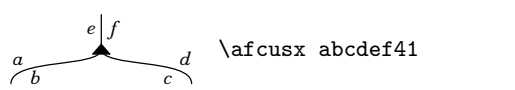
`\afCux abcdef72{Red}`
`{Green}`
`{Blue}`



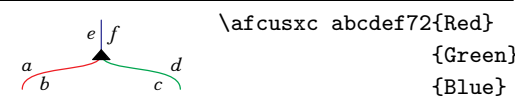
`\afCunx abcd83`



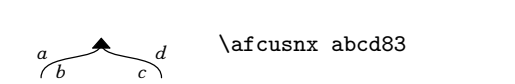
`\afCunxc abcd83{Red}`
`{Green}`



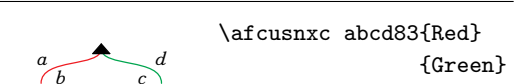
`\afcux abcdef41`



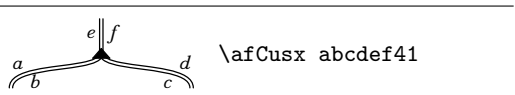
`\afcux abcdef72{Red}`
`{Green}`
`{Blue}`



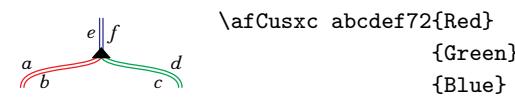
`\afcunx abcd83`



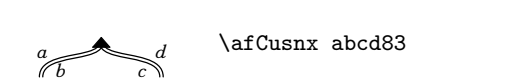
`\afcunxc abcd83{Red}`
`{Green}`



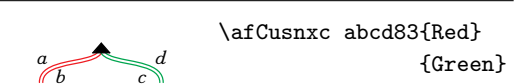
`\afCux abcdef41`



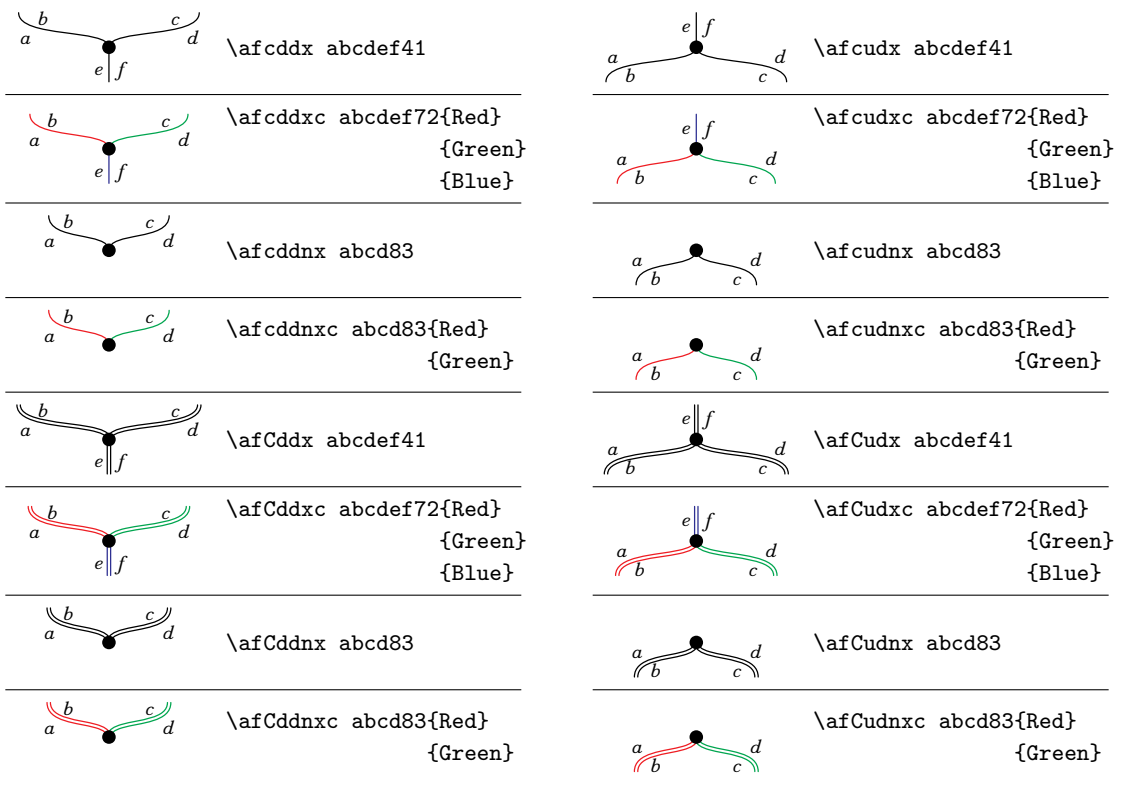
`\afCux abcdef72{Red}`
`{Green}`
`{Blue}`

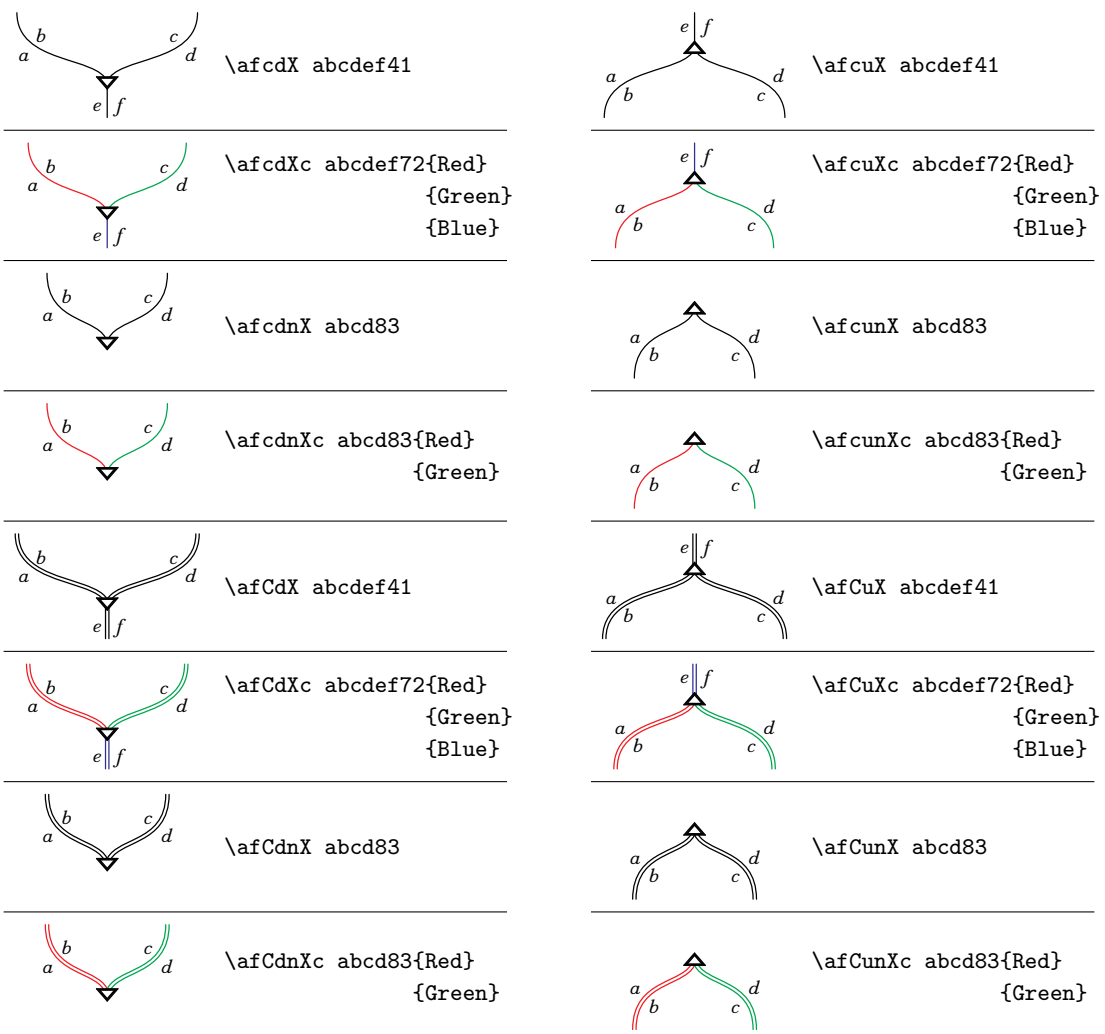


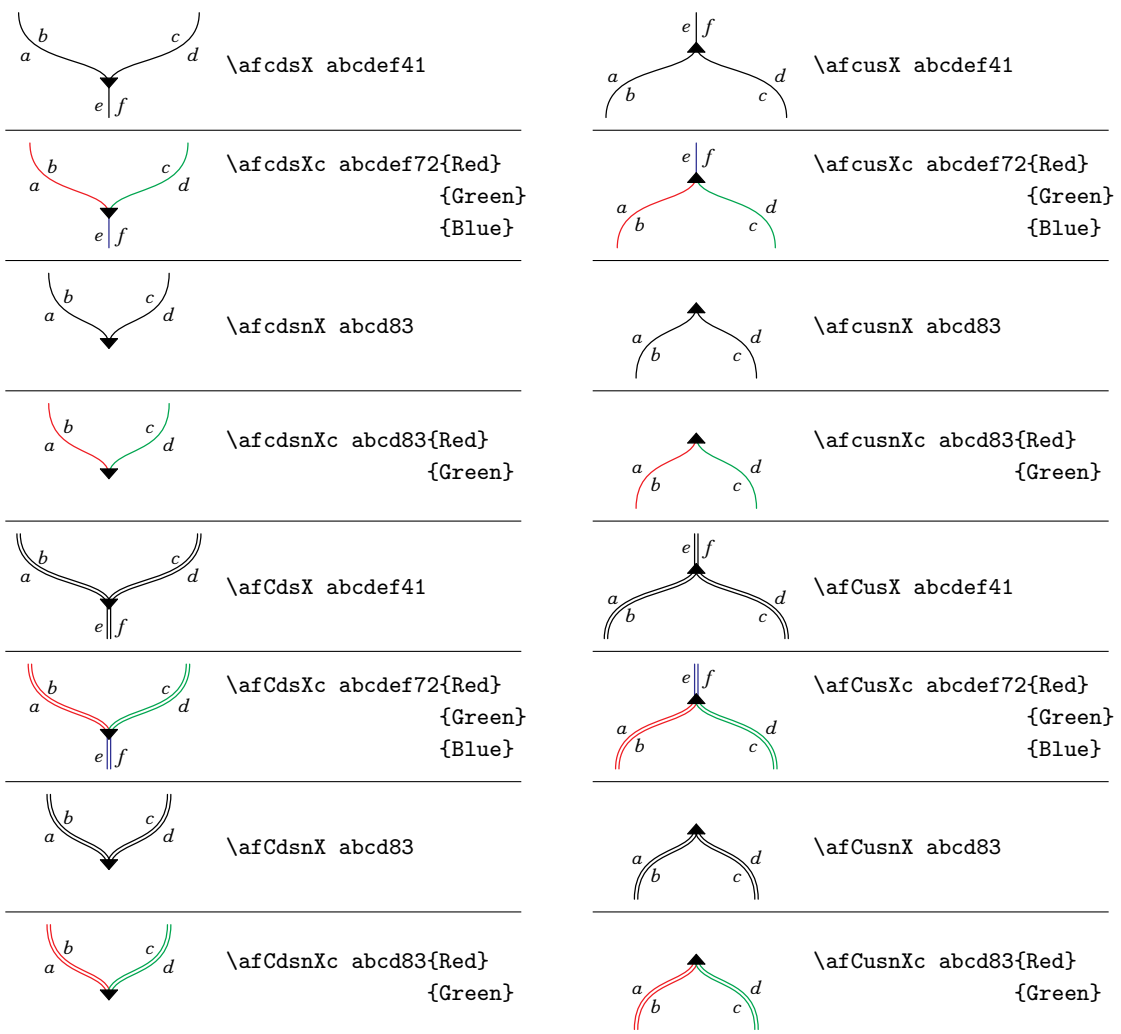
`\afCunx abcd83`

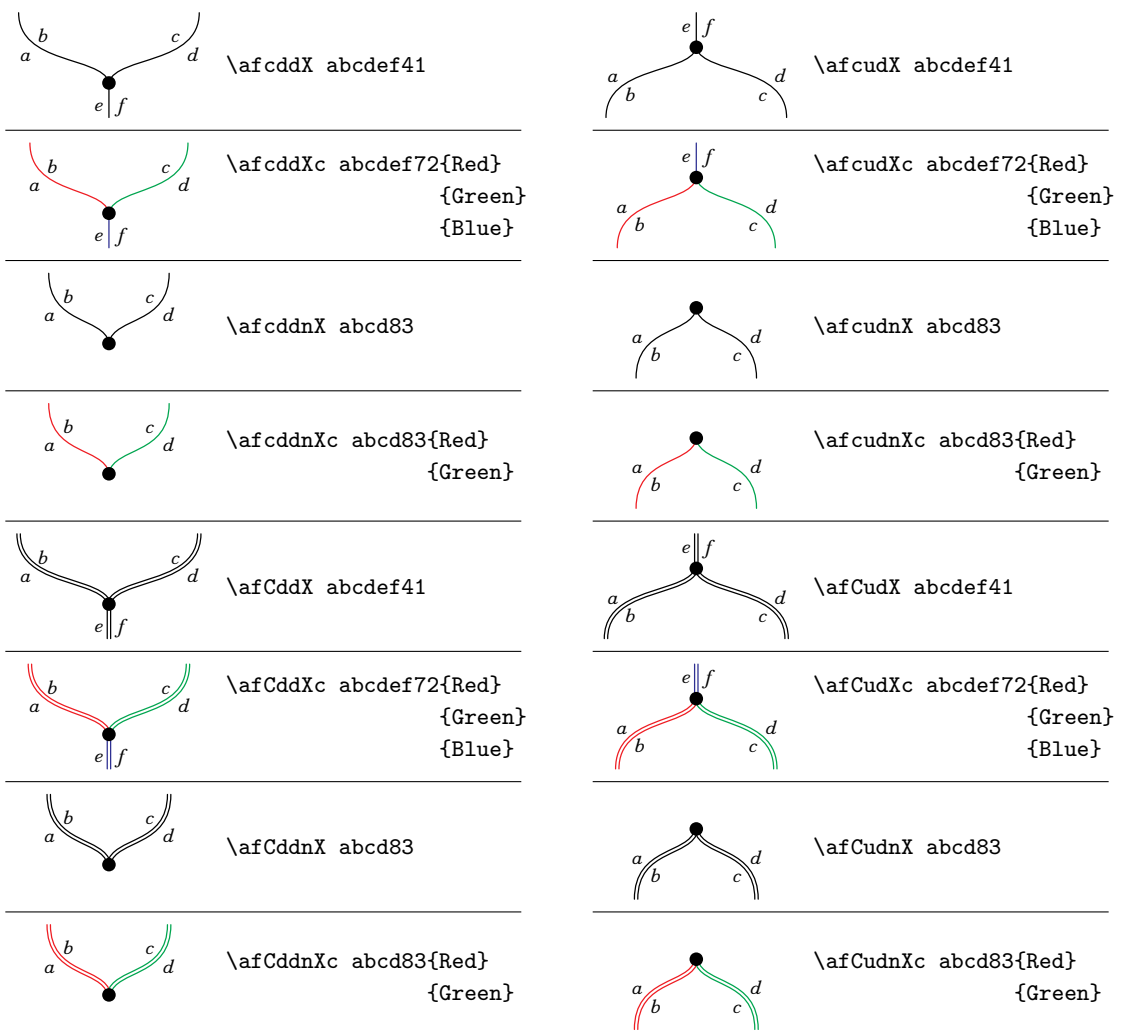


`\afCunxc abcd83{Red}`
`{Green}`



















+ \afbd











⊖ \afbu



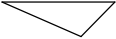



\afv 6
\afvc 6{Red}
$a b$ \afvd 6ab
$a b$ \afvdc 6ab{Red}
$a b$ \afvu 6ab
$a b$ \afvuc 6ab{Red}

\afV 6
\afVc 6{Red}
$a b$ \afVd 6ab
$a b$ \afVdc 6ab{Red}
$a b$ \afVu 6ab
$a b$ \afVuc 6ab{Red}

⋮ \afvdj

 \afjl 44
 \afjlc 44{Red}
 \afjr 44
 \afjrc 44{Red}
 \afcX 48
 \afcxc 48{Red}{Green}
 \afcl 48
 \afclc 48{Red}
 \afcr 48
 \afcrc 48{Red}

 \afJl 44
 \afJlc 44{Red}
 \afJr 44
 \afJrc 44{Red}
 \afCx 48
 \afCxc 48{Red}{Green}
 \afCl 48
 \afClc 48{Red}
 \afCr 48
 \afCrc 48{Red}

 \afex 44
 \afexc 44{Red}{Green}
 \afexpu 542
 \afexpd 825
 \affr 84
 \affrb 48

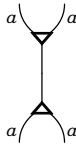
By default the background colour is defined as `\newxycolor{afbackground}{0 0 0 0.12 cmyk}`.

The following dimen parameters are defined:

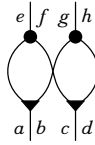
<code>\afelwidth</code>	<code>=425\afunit</code>
<code>\afelheight</code>	<code>=325\afunit</code>
<code>\afthickone</code>	<code>= 40\afunit</code>
<code>\afthickt看</code>	<code>= 40\afunit</code>
<code>\afthickthree</code>	<code>=120\afunit</code>
<code>\afthickfour</code>	<code>= 40\afunit</code>
<code>\aflabeldistance</code>	<code>=220\afunit</code>

The registers govern various parameters in the atomic flow elements (experiment to see which ones). You might be interested especially in changing the last one. The parameters can be globally scaled by invoking `\afsetparams` after having adjusted `\afunit`, which by default is 0.01pt. The parameters `\afthickone` and `\afthickthree` are not effective in pdf \LaTeX .

Each line starting with a coordinate pair (x, y) is an element of the figure. The coordinates are relative, so, for example, you can add or subtract the same vector to all of them and still get the same figure.



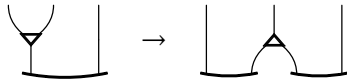
```
\af{
(0,8)*{\afcd a}{a};
(0,0)*{\afcu a}{a}}
```



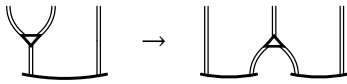
```
\af{
(0,0)*{\afwds ab};
(4,0)*{\afwds cd};
(0,8)*{\afdu ef};
(4,8)*{\afdu gh};
(2,4)*{\afcx 48};
(0,4)*{\afcl 48};
(4,4)*{\afcr 48}}
```



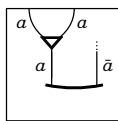
```
\af{
(4,4.5)*{\afvdj};
(0,5 )*{\afcd a}{a};
(2,0 )*{\afiu a}{\bar a}}
```



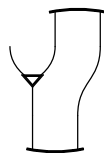
```
\aflower{\af{
( 6, 6)*{\afvd4}{}};
( 3, 0)*{\afiuX}{}}32};
( 0, 4)*{\afcdn}{}}}}
\quad\to\quad
\aflower{\af{
( 6,6)*{\afcu}{}}};
( 0,6)*{\afvd8}{}};
( 12,6)*{\afvd8}{}};
( 10,2)*{\afiu};
( 2,2)*{\afiu}}}}
```



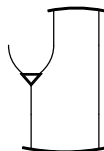
```
\aflower{\af{
( 6, 6)*{\afVd4}{}};
( 3, 0)*{\afIux}{}}32};
( 0, 4)*{\afCdn}{}}}}
\quad\to\quad
\aflower{\af{
( 6,6)*{\afCu}{}}};
( 0,6)*{\afVd8}{}};
( 12,6)*{\afVd8}{}};
( 10,2)*{\afiu};
( 2,2)*{\afiu}}}}
```



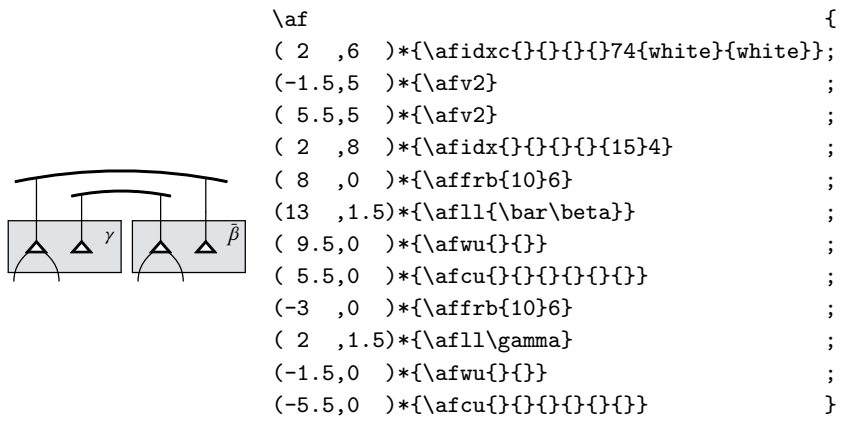
```
\af{
(4,4.5)*{\afvdj};
(0,5 )*{\afcd a}{a};
(1,2.5)*{\affr{10}{13}};
(2,0 )*{\afiu a}{\bar a}}
```



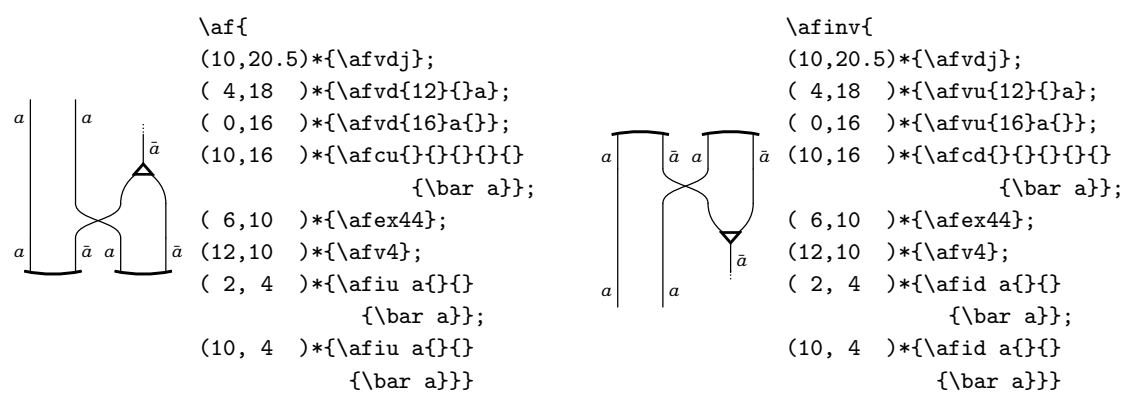
```
\af{
(4,16)*{\afid}{}};
(0, 8)*{\afcd}{}};
(5, 8)*{\afjr28};
(2, 0)*{\afiu}{}}
```



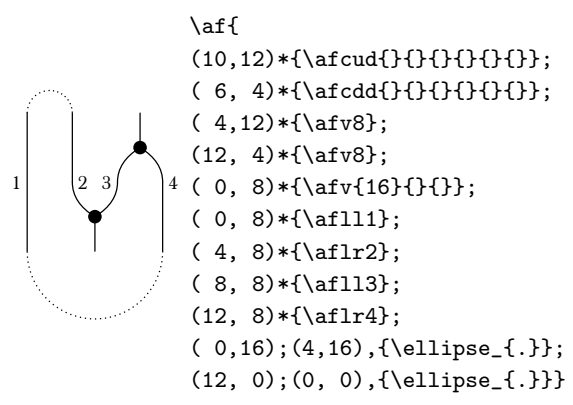
```
{\af{
(4,16)*{\afid}{}};
(0, 8)*{\afcd}{}};
(6, 8)*{\afv8};
(3, 0)*{\afiuX}{}}32}}
```

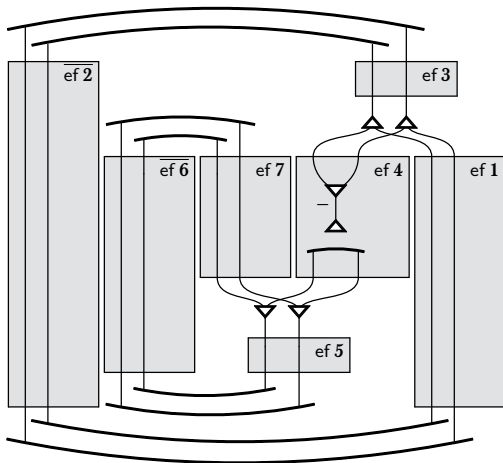


In the next two figures, notice how vertical symmetry is achieved:



The following figure requires `\xyoption{arc}`:



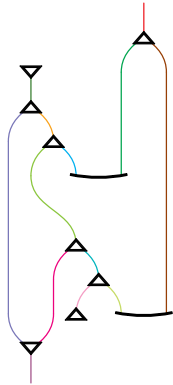


```

\renewcommand{\one }{{\bf\scriptstyle 1}}
\newcommand{\two }{{\bf\scriptstyle 2}}
\newcommand{\three}{{\bf\scriptstyle 3}}
\newcommand{\four }{{\bf\scriptstyle 4}}
\newcommand{\five }{{\bf\scriptstyle 5}}
\newcommand{\six }{{\bf\scriptstyle 6}}
\newcommand{\seven}{{\bf\scriptstyle 7}}
\newcommand{\ef}{{\mathop{\mathsf{ef}}}}

\af
( 1.5 , 4 )*{\affrb{10}{14}} ;
( 0 , 3 )*{\afwu-{} } ;
( 0 , 0 )*{\afidn} ;
( 6.5 , 9.5)*{\afll{\ef\four}} ;
( 0 , 7 )*{\afcdn{}{}{}} ;
(- 2 , - 1.5)*{\afv3} ;
( 2 , - 1.5)*{\afv3} ;
(- 8 , 4 )*{\affrb8{14}} ;
(-10.5 , 3 )*{\afv{12}} ;
(- 8.5 , 4 )*{\afv{14}} ;
(- 4 , 9.5)*{\afll{\ef\seven}} ;
(-11.625,-16 )*{\afiuxc{}{}{}}{10}4
{white}{white} ;
(-11.125,-18 )*{\afiuxc{}{}{}}{59}{16}
{white}{white} ;
(- 3.25 ,-12 )*{\affrb94} ;
(- 6.25 ,-13.5)*{\afv5} ;
(- 3.25 ,-14.5)*{\afv7} ;
( 1.25 ,-11.5)*{\afll{\ef\five}} ;
(- 6.25 ,- 7 )*{\afcdx{}{}{}{}}{17}8 ;
(- 3.25 ,- 7 )*{\afcdx{}{}{}{}}{21}8 ;
(-10.625, 26.8)*{\afidxc{}{}{}}{59}8
{white}{white} ;
(-11.125, 24.8)*{\afidxc{}{}{}}{101}{16}
{white}{white} ;
(- 8.5 ,-22.8)*{\afiuxc{}{}{}}{33}4
{white}{white} ;
(- 8.5 ,-20.5)*{\afiuxc{}{}{}}{59}8
{white}{white} ;
(-16.5 ,- 1.5)*{\affrb8{25}} ;
(-19 ,- 3.5)*{\afv{29}} ;
(-17 ,- 3.5)*{\afv{25}} ;
(-12.5 , 9.5)*{\afll{\overline{\ef\six}}}} ;
(-13.75 , 13 )*{\afidx{}{}{}}{13}8 ;
(-13.75 , 15 )*{\afidx{}{}{}}{21}8 ;
( 11 , -3.5)*{\affrb8{29}} ;
( 8.5 , -4.5)*{\afv{31}} ;
( 10.5 , -5.5)*{\afv{33}} ;
( 15 , 9.5)*{\afll{\ef\one}} ;
( 3.25 , 15 )*{\afcux{}{}{}{}}{21}8 ;
( 6.25 , 15 )*{\afcux{}{}{}{}}{17}8 ;
( 6.25 , 20 )*{\affrb94} ;
( 3.25 , 21 )*{\afv6} ;
( 6.25 , 22 )*{\afv8} ;
( 10.75 , 20.5)*{\afll{\ef\three}} ;
(-25 , 2 )*{\affrb8{40}} ;
(-27.5 , 2 )*{\afv{48}} ;
(-25.5 , 2 )*{\afv{44}} ;
(-21 , 20.5)*{\afll{\overline{\ef\two}}}}

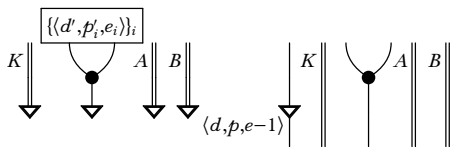
```

```

\af{
( 7,20)*{\afvc{28}{RawSienna}};
(-3, 8)*{\afvc{4}{RubineRed}};
(-2,20)="A";
"A"+( 5, 8)*{\afvc{12}{Green}};
"A"+(-3,10)*{\afcuc{}{}{}{}{}{Periwinkle}
{YellowOrange}
{OliveGreen}};
"A"+(-1, 6)*{\afcunc{}{}{}{}{LimeGreen}
{ProcessBlue}};
"A"+( 3, 2)*{\afiun};
"A"+(-5,-4)*{\afvc{20}{Periwinkle}};
"A"+(-3,14)*{\afwdn};
(-3,18)*{\afjlc48{LimeGreen}};
(-1,14)*{\afcunc{}{}{}{}{RubineRed}
{Aquamarine}};
( 1,10)*{\afcunc{}{}{}{}{Lavender}
{SpringGreen}};
( 5, 6)*{\afiun};
(-1, 6)*{\afwun};
( 5,38)*{\afcuc{}{}{}{}{}{Green}
{RawSienna}
{Red}};
(-5, 2)*{\afcvc{}{}{}{}{}{Periwinkle}
{RubineRed}
{DarkOrchid}}

```



```

\af{
(-1 , 4 )*{\afAd{}{}}};
(-1 ,10 )*{\afVd4K{}}};
( 4.5,14 )*{\affr94}}};
( 4.5,14 )*{\strut\scriptstyle
{\langle d',p'_i,e_i\rangle\}_i};
( 4.5, 8 )*{\afcdd{}{}{}{}{}}};
( 4.5, 4 )*{\afadn}}};
(10 , 4 )*{\afAd{}{}}};
(10 ,10 )*{\afVd4A{}}};
(13 , 4 )*{\afAd{}{}}};
(13 ,10 )*{\afVd4B{}}};
(22 , 2 )*{\afv4}}};
(22 , 4 )*{\afad{}{}}};
(22 ,10 )*{\afv4}}};
(22 , 2.5)*{\afll{\strut\scriptstyle
\langle d,p,e-1\rangle}}};
(25 , 6 )*{\afVd{12}K{}}};
(29 , 2 )*{\afv4}}};
(29 , 8 )*{\afcdd{}{}{}{}{}}};
(33 , 6 )*{\afVd{12}A{}}};
(36 , 6 )*{\afVd{12}B{}}};

```


7. SMASH MACROS

Compare:

α aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ β α $\rho \parallel S$ aaa aaa aaa aaa aaa aaa aaa aaa β α aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ β	aaa aaa aaa aaa aaa aaa aaa aaa $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$ $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$ aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$ <hr/> α aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ $\$ \backslash \text{vldownsmash}$ β $\{\backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\}\$$ $\rho \parallel S$ $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$ β aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ β $\$ \backslash \text{vlupsmash}\{\backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\}\$$ <hr/> α aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$ β $\$ \backslash \text{vlsmash}\{\backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\}\$$ $\rho \parallel S$ aaa aaa aaa aaa aaa aaa aaa aaa β aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\rho \parallel S$ β $\$ \backslash \text{vlder}\{\rho\}\{\text{cal } S\}\{\beta\}\{\alpha\}\$$
---	---

8. ACKNOWLEDGEMENTS

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