

VIRGINIA LAKE L^AT_EX MACROS V3.3

<http://alessio.guglielmi.name/res/vl>

Alessio Guglielmi

University of Bath

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To use this package: `\usepackage{virginialake}`.

If you don't need any graphical elements (most of which have curved lines), you can save some processing time and L^AT_EX resources by doing `\usepackage[noxy]{virginialake}`. In this case, the package does not load XY-pic.

For elements with curved lines there are two options:

- compile the L^AT_EX file with `latex + dvips`: this takes longer but gives a slightly better output;
- compile with `pdflatex`: in this case, some graphical elements will have lower quality, but compilation is faster.

Calling the package with `\usepackage[goodsyntax]{virginialake}` produces formulae or structures in the deep-inference style.

Calling the package with `\usepackage[lutzsyntax]{virginialake}` produces formulae or structures in the traditional style, except that it uses square brackets for disjunctions.

To compile this manual from the source make sure that `\write18` is enabled; this can be achieved by invoking `pdflatex` with the `--shell-escape` option.

If T_EX runs out of resources while using Virginia Lake, try invoking `\usepackage{etex}` immediately after `\documentclass`.

For use with Beamer Inside the `frame` environment in beamer parentheses and punctuation marks have a special treatment and can cause problems, notably when using square brackets for optional arguments or round brackets for atomic flow coordinates. The solution is to use the option `fragile` for every frame where Virginia Lake macros are used (as in `\begin{frame}[fragile]`). If a frame contains only one slide the option `fragile=singleslide` insures greater speed.

1. SYMBOLS

The following symbols are defined:

1	<code>\one</code>	c↑	<code>\cU</code>
0	<code>\zer</code>	w↓	<code>\wD</code>
f	<code>\fff</code>	w↑	<code>\wU</code>
t	<code>\ttt</code>	q↓	<code>\qD</code>
s	<code>\sw</code>	q↑	<code>\qU</code>
m	<code>\me</code>	BV	<code>\BV</code>
i↓	<code>\iD</code>	SBV	<code>\SBV</code>
i↑	<code>\iU</code>	KS	<code>\KS</code>
c↓	<code>\cD</code>	SKS	<code>\SKS</code>

2. FORMULAE AND STRUCTURES

Normal formulae:

$$\begin{array}{ll} ((A \vee (B \wedge C) \rightarrow \neg D) \leftrightarrow E) \leftarrow F & ((A \text{\vlor}(B \text{\vlan} C) \text{\vlim}\text{\vlne} D) \text{\vldi} E) \text{\vlmi} F \\ (A \supset (B \subset C)) & (A \text{\vljm} (B \text{\vlmj} C)) \\ (A \wp (B \otimes (C \triangleleft D))) \multimap E & (A \text{\vlp} (B \text{\vlt} (C \text{\vlse} D))) \text{\vlli} E \end{array}$$

The command `\vllineartrue` triggers linear logic additive conjunction and disjunction. It is reversed by `\vllinearfal`, which is the default:

$$((A \oplus (B \& C) \rightarrow \neg D) \leftrightarrow E) \leftarrow F \quad \text{\vllineartrue} \\ ((A \text{\vlor}(B \text{\vlan} C) \text{\vlim}\text{\vlne} D) \text{\vldi} E) \text{\vlmi} F$$

Logical operators become smaller than usual if `\vsmallops` is issued. The macro `\vbin` creates a logical operator: *e.g.*, `\vlse` is `\vbin\triangleleft`.

Note that the new macros `\lt` and `\gt` are defined, respectively, for `<` and `>` (whose corresponding ASCII characters have special meaning when dealing with formulae).

3. FORMULAE AND STRUCTURES VIA PUNCTUATION AND PARENTHESES

For typesetting formulae, it is possible to use punctuation marks and parentheses instead of macros, as shown below.

The following happens if the package is called without the option `[goodsyntax]`, or if the command `\vlnogoodsyntax` is issued:

$$\begin{array}{ll} A \vee \bar{B} & \{\text{\vls}[A.\neg B]\} \\ A \vee \bar{B} \vee \dots \vee D \vee E & \{\text{\vls}[A.\neg B.\text{\vldots}.D.E]\} \\ A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G & \{\text{\vls}[A..(\neg B.[C;\lt D;(E;F)\gt]).G]\} \\ (A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G) & \{\text{\vlsbr}[A.(\neg B..[C;\lt D;(E;F)\gt]).G]\} \\ \{A \vee (\bar{B} \wedge (C \wp (D \triangleleft (E \otimes F)))) \vee G\} & \{\text{\vlscn}[A.(\neg B.[C;\lt D;;(E;F)\gt]).G]\} \\ A\{\} & \text{\vlhole} \end{array}$$

The following happens if the package is called with the option `[goodsyntax]`, or if the command `\vlgoodsyntax` is issued:

$$\begin{array}{ll} [A \bar{B}] & \{\text{\vls}[A.\neg B]\} \\ [A \bar{B} \dots D E] & \{\text{\vls}[A.\neg B.\text{\vldots}.D.E]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\text{\vls}[A..(\neg B.[C;\lt D;(E;F)\gt]).G]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\text{\vlsbr}[A.(\neg B..[C;\lt D;(E;F)\gt]).G]\} \\ [A (\bar{B} [C \langle D (E F) \rangle]) G] & \{\text{\vlscn}[A.(\neg B.[C;\lt D;;(E;F)\gt]).G]\} \\ A\{\} & \text{\vlhole} \end{array}$$

The following happens if the package is called with the option `[lutzsyntax]`, or if the command `\vllutzsyntax` is issued:

$$\begin{array}{ll} A \vee \bar{B} & \{\text{\vls}[A.\neg B]\} \\ A \vee \bar{B} \vee \dots \vee D \vee E & \{\text{\vls}[A.\neg B.\text{\vldots}.D.E]\} \\ A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G & \{\text{\vls}[A..(\neg B.[C;\lt D;(E;F)\gt]).G]\} \\ [A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G] & \{\text{\vlsbr}[A.(\neg B..[C;\lt D;(E;F)\gt]).G]\} \\ \{A \vee (\bar{B} \wedge [C \wp \langle D \triangleleft (E \otimes F) \rangle]) \vee G\} & \{\text{\vlscn}[A.(\neg B.[C;\lt D;;(E;F)\gt]).G]\} \\ A\{\} & \text{\vlhole} \end{array}$$

The command `\vlsmallbrackets` corrects a possible problem with the size of brackets:

$$\begin{array}{c} (a \vee b) \wedge (\bar{a} \vee \bar{b}) \quad \{\backslash vls([a.b].[a..b])\} \\ \hline (a \vee b) \wedge (\bar{a} \vee \bar{b}) \quad \{\backslash vlsmallbrackets \\ \quad \backslash vls([a.b].[a..b])\} \end{array}$$

This also holds for the alternative syntax style:

$$\begin{array}{c} ([a\,b]\,[\bar{a}\,\bar{b}]) \quad \{\backslash vls([a.b].[a..b])\} \\ \hline ([a\,b]\,[\bar{a}\,\bar{b}]) \quad \{\backslash vlsmallbrackets \\ \quad \backslash vls([a.b].[a..b])\} \end{array}$$

The command `\vlnosmallbrackets` undoes `\vlsmallbrackets`.

The `\vls` macro works by redefining ‘:, ‘[’, ‘]’, ‘(’ and ‘)’. This, of course, can cause several problems. In practice, they are rare, except when `\vls` is an argument of a macro. In this case, one needs to ‘update’ the macro by calling `\vlupdate`. Compare

$$[a.b] \quad \begin{array}{l} \text{\newcommand{\vltest}[1]{#1}} \\ \text{\vltest{\vls[a.b]}} \end{array}$$

with

$$a \vee b \quad \begin{array}{l} \text{\newcommand{\vltest}[1]{#1}} \\ \text{\vlupdate\vltest} \\ \text{\vltest{\vls[a.b]}} \end{array}$$

This mechanism, in principle, works with every macro.

Sometimes it is not possible to use `\vlupdate`. For example, in the following situation (requiring the `amsmath` package):

$$\begin{aligned} \beta_k &= \bigwedge_{k \leq i \leq n} [-c_i. - d_i] \quad \text{for } 1 < k \leq n, & \backslash begin{align*} \\ && \backslash beta_k \&=\& \backslash bigwedge_{k \leq i \leq n} \\ && \quad \backslash vls[-c_i.-d_i] \\ && \quad \& \\ && \quad \backslash text{for } \$1 < k \leq n \$, \\ \gamma_i &= (\beta_{i+1}.c_i) \quad \text{for } 1 \leq i < n, & \backslash gamma_i \&=\& \backslash vls(\backslash beta_{i+1}.c_i) \& \\ \delta_i &= (\beta_{i+1}.d_i) \quad \text{for } 1 \leq i < n, & \backslash text{for } \$1 \leq i < n \$, \\ \gamma_n &= c_n, & \backslash delta_i \&=\& \backslash vls(\backslash beta_{i+1}.d_i) \& \\ \delta_n &= d_n. & \backslash text{for } \$1 \leq i < n \$, \\ && \backslash gamma_n \&=\& c_n, \\ && \backslash delta_n \&=\& d_n. \\ && \backslash end{align*} \end{aligned}$$

The solution is to use the macros `\vlstore`, `\vlread` and `\lt` as follows:

$$\begin{aligned} \beta_k &= \bigwedge_{k \leq i \leq n} \bar{c}_i \vee \bar{d}_i \quad \text{for } 1 < k \leq n, & \backslash begin{align*} \\ && \backslash vlstore \% \\ && \backslash beta_k \&=\& \backslash bigwedge_{k \leq i \leq n} \\ && \quad \backslash vls[-c_i.-d_i] \\ && \quad \& \\ && \quad \backslash text{for } \$1 \lt k \leq n \$, \\ \gamma_i &= \beta_{i+1} \wedge c_i \quad \text{for } 1 \leq i < n, & \backslash gamma_i \&=\& \backslash vls(\backslash beta_{i+1}.c_i) \& \\ \delta_i &= \beta_{i+1} \wedge d_i \quad \text{for } 1 \leq i < n, & \backslash text{for } \$1 \leq i < n \$, \\ \gamma_n &= c_n, & \backslash delta_i \&=\& \backslash vls(\backslash beta_{i+1}.d_i) \& \\ \delta_n &= d_n. & \backslash text{for } \$1 \leq i < n \$, \\ && \backslash gamma_n \&=\& c_n, \\ && \backslash delta_n \&=\& d_n \backslash vldot \\ && \} \\ && \backslash begin{align*} \\ && \backslash vlread \\ && \backslash end{align*} \end{aligned}$$

Note the use of `\vldot` in the place of `.` in the last line of the display. The commands `\vldot` for `.` and `\vlsqb` for `[` and `\vlsqbr` for `]` and `\vlrobr` for `(` and `\vlrobrr` for `)` are provided.

There's the command `\vlnos`, which sets `'`, `[`, `]`, `(` and `)` to their normal behaviour, for example:

```
a ∧ (. [ b ∧ c ].) {\vls(a.{\vlnos (. [{\vls(b.c)}] .))})}
```

4. DEEP INFERENCE AND SEQUENT CALCULUS DERIVATIONS

The following macros, by default, accept the commands for formulae and structures, as shown before. However, this might create problems in some circumstances, because the characters `'`, `[`, `]`, `(` and `)` are defined in a special way. To avoid such problems, in case the commands for formulae/structures are not wanted, you can use the command `\vlnostructuresyntax`: every successive invocation of the commands for derivations does not redefine any character.

$\rho \frac{A}{B} \star \text{\vlderivation} \{ \text{\vlin}\{\rho\}\{\star\}\{B\} \{ \text{\vlhy} \{A\} \} \}$	$\rho \frac{A \quad B}{C} \star \text{\vlderivation} \{ \text{\vliin}\{\rho\}\{\star\}\{C\} \{ \text{\vlhy} \{A\} \} \{ \text{\vlhy} \{B\} \} \}$
$\rho \frac{A}{B} \star \text{\vlderivation} \{ \text{\vlid}\{\rho\}\{\star\}\{B\} \{ \text{\vlhy} \{A\} \} \}$	$\rho \frac{A \quad B}{C} \star \text{\vlderivation} \{ \text{\vliid}\{\rho\}\{\star\}\{C\} \{ \text{\vlhy} \{A\} \} \{ \text{\vlhy} \{B\} \} \}$
$\rho \frac{A}{B} \star \text{\vlderivation} \{ \text{\vliq}\{\rho\}\{\star\}\{B\} \{ \text{\vlhy} \{A\} \} \}$	$\rho \frac{A \quad B}{C} \star \text{\vlderivation} \{ \text{\vliiq}\{\rho\}\{\star\}\{C\} \{ \text{\vlhy} \{A\} \} \{ \text{\vlhy} \{B\} \} \}$
$\rho \frac{A}{B} \star \text{\vlinf}\{\rho\}\{\star\}\{B\}\{A\}$	$\rho \frac{A \quad B}{C} \star \text{\vliinf}\{\rho\}\{\star\}\{C\}\{A\}\{B\}$
$\rho \frac{A}{B} \star \text{\vlidf}\{\rho\}\{\star\}\{B\}\{A\}$	$\rho \frac{A \quad B}{C} \star \text{\vliidf}\{\rho\}\{\star\}\{C\}\{A\}\{B\}$
$\rho \frac{A}{B} \star \text{\vliqf}\{\rho\}\{\star\}\{B\}\{A\}$	$\rho \frac{A \quad B}{C} \star \text{\vliiqf}\{\rho\}\{\star\}\{C\}\{A\}\{B\}$

$\rho \frac{A \quad B \quad C}{D} \star$	<pre>\vlderivation { \vliiin{\rho}{\star}{D} { \vlhy {A} { \vlhy {B} { \vlhy {C}}}}</pre>
$\rho \frac{A \quad B \quad C}{D} \star$	<pre>\vlderivation { \vliiid{\rho}{\star}{D} { \vlhy {A} { \vlhy {B} { \vlhy {C}}}}</pre>
$\rho \frac{A \quad B \quad C}{\overline{D}} \star$	<pre>\vlderivation { \vliiqf{\rho}{\star}{D} { \vlhy {A} { \vlhy {B} { \vlhy {C}}}}</pre>
$\rho \frac{A \quad B \quad C}{D} \star$	<pre>\vlderivation { \vliiinf{\rho}{\star}{D}{A}{B}{C}}</pre>
$\rho \frac{A \quad B \quad C}{D} \star$	<pre>\vlderivation { \vliiidf{\rho}{\star}{D}{A}{B}{C}}</pre>
$\rho \frac{A \quad B \quad C}{\overline{D}} \star$	<pre>\vlderivation { \vliiqf{\rho}{\star}{D}{A}{B}{C}}</pre>

$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiin{\rho}{\star}{E} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D}}})}</pre>
$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiid{\rho}{\star}{E} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D}}})}</pre>
$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiinq{\rho}{\star}{E} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D}}})}</pre>
$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiinf{\rho}{\star}{E}{A}{B}{C}{D}</pre>
$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiidf{\rho}{\star}{E}{A}{B}{C}{D}</pre>
$\rho \frac{A \quad B \quad C \quad D}{E} \star$	<pre>\vlderivation { \vliiinqf{\rho}{\star}{E}{A}{B}{C}{D}</pre>

$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vlderivation { \vliiiin{\rho}{\star}{F} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} } } } } } }</pre>
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vlderivation { \vliiid{\rho}{\star}{F} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} } } } } } }</pre>
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vlderivation { \vliiiiq{\rho}{\star}{F} { \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} } } } } } }</pre>
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vlderivation { \vliiiinf{\rho}{\star}{F}{A}{B}{C}{D}{E} }</pre>
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vliiiidf{\rho}{\star}{F}{A}{B}{C}{D}{E}</pre>
$\rho \frac{A \ B \ C \ D \ E}{F} \star$	<pre>\vliiiiqf{\rho}{\star}{F}{A}{B}{C}{D}{E}</pre>

$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliini{\rho}{\star}{G}{ \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} { \vlhy {F}}})}}</pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliidi{\rho}{\star}{G}{ \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} { \vlhy {F}}})}}</pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliiq{\rho}{\star}{G}{ \vlhy {A} { \vlhy {B} { \vlhy {C} { \vlhy {D} { \vlhy {E} { \vlhy {F}}})}}</pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliinf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F}}</pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliidf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F}}</pre>
$\rho \frac{A \ B \ C \ D \ E \ F}{G} \star$	<pre>\vlderivation { \vliliiqf{\rho}{\star}{G}{A}{B}{C}{D}{E}{F}}</pre>

The commands `\vlsmallleftlabels` and `\vlnosmallleftlabels` control the size of labels at the left of $\frac{\quad}{\quad}$ (the ρ s above); the default is small size.

$\Pi \parallel_S$	$\backslash vlderivation$	{
$B \vee C$	$\backslash vlp{\{ \backslash Pi \} \{ \backslash cal S \} \{ \backslash vls[B.C] \}}$	}
<hr/>		
$\Pi \parallel_S$	$\backslash vlderivation$	{
$B \vee C$	$\backslash vlpd{\{ \backslash Pi \} \{ \backslash cal S \} \{ \backslash vls[B.C] \}}$	}
<hr/>		
$\Pi \parallel_S$	$\backslash vlderivation$	{
$B \vee C$	$\backslash vlpf{\{ \backslash Pi \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{ \backslash the \backslash toks0 \}}$	}
<hr/>		
$\Pi \parallel_S$	$\backslash vlderivation$	{
$B \vee C$	$\backslash vlp{\{ \backslash Pi \} \{ \backslash cal S \} \{ \backslash vls[B.C] \}}$	}
<hr/>		
A	$\backslash vlderivation$	{
$\Delta \parallel_S$	$\backslash vld{\{ \Delta \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{$	
$B \vee C$	$\backslash vly$	{A }}}
<hr/>		
A	$\backslash vlderivation$	{
$\Delta \parallel_S$	$\backslash vldf{\{ \Delta \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{$	
$B \vee C$	$\backslash vly$	{A }}}
<hr/>		
A	$\backslash vlderivation$	{
$\Delta \parallel_S$	$\backslash vldf{\{ \Delta \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{$	
$B \vee C$	$\backslash vly$	{A }{ \backslash the \backslash toks0 } }
<hr/>		
A	$\backslash vlderivation$	{
$\Delta \parallel_S$	$\backslash vld{\{ \Delta \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{ A }$	
<hr/>		
A	$\backslash vlderivation$	{
$\Delta \parallel_S$	$\backslash vld{\{ \Delta \} \{ \backslash cal S \} \{ \backslash vls[B.C] \} \{ A }$	
$B \vee C$		}

$\Pi \prod_S$	<code>\vlderivation {</code>
A	<code>\vlin{\rho}{\rm note}{\vls[B.C]}{}</code>
$\rho \frac{A}{B \vee C} \text{note}$	<code>\vlpr{\Pi}{\cal S}{A}}</code>
<hr/>	<hr/>
A	<code>\vlderivation {</code>
$\Delta \prod_S$	<code>\vlin{\rho}{ }{C}{}</code>
B	<code>\vldf{\Delta}{\cal S}{B}{}</code>
$\rho \frac{A}{C}$	<code>\vlhy {A}}}}</code>
<hr/>	<hr/>
A	<code>\vlderivation {</code>
$\Delta \prod_{S''}$	<code>\vlin{\rho}{ }{D}{}</code>
B	<code>\vldf{\Delta'}{\cal S}{D}{}</code>
$\Delta \prod_{S'}$	<code>\vldf{\Delta'}{\cal S'}{C}{}</code>
C	<code>\vldf{\Delta}{\cal S''}{B}{}</code>
$\Delta'' \prod_S$	<code>\vlhy {A}}}}</code>
D	
<hr/>	<hr/>
A	<code>\vlderivation {</code>
$\rho \frac{A}{B}$	<code>\vldf{\Delta'}{}{D}{}</code>
$\Delta \prod$	<code>\vldf{\Delta}{ }{C}{}</code>
C	<code>\vlin{\rho}{ }{B}{}</code>
$\Delta \prod$	<code>\vlhy {A}}}}</code>
D	
<hr/>	<hr/>
A	<code>\vlderivation {</code>
$\rho \frac{A}{B}$	<code>\vlin{\rho'}{\quad}{D}{}</code>
$\rho' \frac{A}{\overline{C}}$	<code>\vliq{\rho'}{C}{}</code>
$\rho'' \frac{A}{D}$	<code>\vlid{\rho}{B}{}</code>
.	<code>\vlhy {A}}}}</code>
<hr/>	<hr/>
$A \wedge A'$	<code>\vlderivation {</code>
$\rho \frac{A \wedge A'}{B \wedge B'}$	<code>\vliiq{\rho}{\star}{\vls(E.E')}{}</code>
$\Delta \prod$	<code>\vldf{\Delta}{\vls(D.D')}{}</code>
$C \wedge C'$	<code>\vldf{\Delta}{\vls(C.C')}{}</code>
$\Delta' \prod$	<code>\vlin{\rho}{\vls(B.B')}{}</code>
$D \wedge D'$	<code>\vlhy {\vls(A.A')}}}</code>
$\rho \frac{A \wedge A' \quad \Delta \prod \quad C \wedge C' \quad \Delta' \prod \quad D \wedge D'}{E \wedge E'} \star$	<code>{</code>
	<code>\vlin{\rho'}{\star}{\vls(D.D')}{}</code>
	<code>\vlin{\rho'}{\vls(C.C')}{}</code>
	<code>\vlin{\rho}{\vls(B.B')}{}</code>
	<code>\vlhy {\vls(A.A')}}}</code>

Note `\vldots` in

$$\begin{array}{l}
 \gamma \\
 \parallel \\
 \alpha_1 \\
 \tau \frac{\alpha_1}{\alpha_1 \sigma_1} \vlderivation {
 \vldf{}{}{\beta} \\
 \vlin{\rho}{}{\alpha_n \sigma_n} \\
 \vdots \\
 \vldf{}{}{\vldots} \\
 \parallel \\
 \vlin{\tau}{}{\alpha_1 \sigma_1} \\
 \rho \frac{\alpha_n}{\alpha_n \sigma_n} \vlderivation {
 \vldf{}{}{\alpha_1} \\
 \vlhy {\gamma} \\
 \parallel \\
 \beta
 }
 }
 \end{array}$$

The command \vlx provides some extra space, as in

```

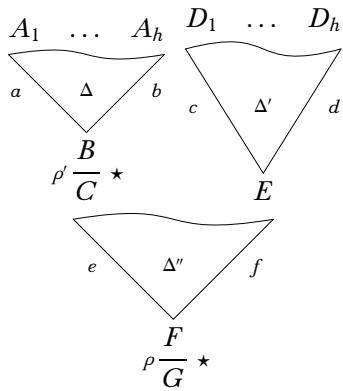
\vlderivation
\vldd{\Psi'}{}{\beta}
\vldd{\Psi} {}{\xi{\left( \begin{array}{c} \alpha \\ \Psi \\ \left( \begin{array}{c} t \\ \overline{a \vee \bar{a}} \wedge (\bar{a} \vee \bar{\bar{a}}) \end{array} \right) \\ \left( \begin{array}{c} a \\ \overline{a \wedge a} \wedge (\bar{a} \vee \bar{\bar{a}}) \\ \overline{a \wedge ((a \wedge \bar{a}) \vee \bar{a})} \vee \bar{a} \\ \overline{a \wedge \bar{a}} \vee \overline{a \wedge \bar{\bar{a}}} \\ f \end{array} \right) \\ \Psi \\ \beta \end{array} \right)}}
\vlinf{}{}{\vls[\vlinf{}]}
\vlinf{}{}{\vlsbr[\vlderivation]
\vlinf{}{}{\vls[\vlinf{}]}
\vlinf{}{}{\vls(a.-a)}
\vlinf{}{}{\vls(a.-a)}
\vlhy{\vls(\vlinf{}{}{\vls(a.[(a.-a).-a])})}}
\vlhy{\vls(\vlinf{}{})}
\vlhy{\vls(a.a)}
\vlhy{a}
\vlhy{[-a. -a]}
)
\vlx.\vlx
{-a}
]
\vlhy{\vls(\vlinf{}{}{\vls[a.-a])t.[ -a.-a])}
\vlhy{\vls(\vlinf{}{}{\vls[a.-a])t.[ -a.-a])}
\vlhy{\right\}}
\vlhy{\alpha}
\right\}

```

The following use XY-pic (see at the beginning of this manual for instructions).

	<pre>\vlderivation { \vltr{\Delta}{B }{ \vlhy {A_1 } { \vlhy {\dots} { \vlhy {A_h } } } } }</pre>
	<pre>\vltreeder{\Delta}{B}{A_1}{\dots}{A_h}</pre>
	<pre>\vlderivation { \vltrl{\Delta}{a}{b}{B }{ \vlhy {A_1 } { \vlhy {\dots} { \vlhy {A_h } } } } }</pre>
	<pre>\vltreederl{\Delta}{a}{b}{B}{A_1}{\dots}{A_h}</pre>
	<pre>{\vlnostructuressyntax \vlderivation { \vltrf{\Delta}{B }{ \vlhy {A_1 } { \vlhy {\dots} { \vlhy {A_h } } } } }</pre>
	<pre>\toks0={0.5} \vlderivation { \vltrf{\Delta}{\vls{B,C}}{ \vlhy {A_1 } { \vlhy {\dots} { \vlhy {A_h } } } } }</pre>
	<pre>\toks0={2.5} \vlderivation { \vltrlf{\Delta}{a}{b}{\vls{B,C}}{ \vlhy {A_1 } { \vlhy {\dots} { \vlhy {A_h } } } } }</pre>

	<pre>\dimen0=-1pc \vllderivation \vlinf{\rho}{\star}{G} { \vltr{\Delta'}{F} { \vlinf{\rho'}{\star}{C} { \vltr{\Delta}{B} { \vlhy {A_1} } } \vlhy {\dots} } } \vlhy {A_h} } } \vlhy {\dots} } } \vlhy {\kern\dimen0} } } \vltr{\Delta'} {E} { \vlhy {D_1} } } \vlhy {\dots} } } \vlhy {\kern\dimen0} } } \vltr{\Delta''}{F} { \vlhy {G} } } \vlhy {\star} } } \vllderivation \vliiin{\rho}{\star}{G} { \vltr{\Delta}{B} { \vlhy {A_1} } } \vlhy {A_2} } } \vlhy {A_3} } } \vltr{\Delta'} {D} { \vlhy {C_1} } } \vlhy {C_2} } } \vlhy {C_3} } } \vltr{\Delta''}{F} { \vlhy {E_1} } } \vlhy {E_2} } } \vlhy {E_3} } } \vltr{\Delta'} {F} { \vlhy {D} } } \vlhy {C_1} } } \vlhy {C_2} } } \vlhy {C_3} } } \vltr{\Delta''}{F} { \vlhy {E_1} } } \vlhy {E_2} } } \vlhy {E_3} } } \vltr{\Delta'} {F} { \vlhy {E_1} } } \vlhy {E_2} } } \vlhy {E_3} } } }</pre>



```

{\vlnostructuressyntax
\vlderivation
\vlin {\rho} {\star}{G}
\vltrl {\Delta''}{e}{f} {F}
\vlin {\rho'} {\star}{C}
\vltrl {\Delta}{a}{b} {B}
\vlhy {A_1}
\vlhy {\dots}
\vlhy {A_h}
\vlhy {B}
\vlhy {C}
\vlhy {\kern-1pc}
\vlhy {\Delta'}
\vlhy {E}
\vlhy {\dots}
\vlhy {A_h}
\vlhy {\kern-1pc}
\vlhy {\Delta''}
\vlhy {F}
\vlhy {G}
\vlhy {\dots}
\vlhy {D_h}
\vlhy {\kern-1pc}
\vlhy {\{1.6\}}}

```

5. OPEN DEDUCTION

The following macros for open deduction do not reverse the order of formulae.

$\rho \frac{A}{B \vee C} \star$	$\{\{A\}\}$	$\backslash od\{\backslash odif\{\backslash odh$
	$\{\{\backslash rho\}\{[B.C]\}\{\backslash star\}\}$	$\}$
$\rho \frac{A}{B \vee C} \star$	$\{\{A\}\}$	$\backslash od\{\backslash odof\{\backslash odh$
	$\{\{\backslash rho\}\{[B.C]\}\{\backslash star\}\}$	$\}$
$\rho \frac{A}{B \vee C} \star$	$\{\{A\}\}$	$\backslash od\{\backslash odI\{\backslash odh$
	$\{\{\backslash rho\}\{[B.C]\}\{\backslash star\}\}$	$\}$
$\rho \frac{A}{B \vee C} \star$	$\backslash odn\{A\}\{\backslash rho\}\{[B.C]\}\{\backslash star\}$	
$\rho \frac{A}{B \vee C} \star$	$\backslash odn\{A\}\{\backslash rho\}\{[B.C]\}\{\backslash star\}$	
$\rho \frac{A}{B \vee C} \star$	$\backslash odN\{A\}\{\backslash rho\}\{[B.C]\}\{\backslash star\}$	
$\rho \frac{\Pi \parallel S}{B \vee C} \star$	$\{\{\Pi\}\{[B.C]\}\{\backslash cal S\}\}$	$\backslash od\{\backslash odp$
$\rho \frac{\Pi \parallel S}{B \vee C} \star$	$\{\{\Pi\}\{[B.C]\}\{\backslash cal S\}\}$	$\backslash od\{\backslash odP$
$\rho \frac{\Pi \parallel S}{B \vee C} \star$	$\{\{\Pi\}\{[B.C]\}\{\backslash cal S\}\}$	$\backslash od\{\backslash odpx$
		$\{\{\backslash Pi\}\{[B.C]\}\{\backslash cal S\}\{\the\toks0\}\}$
$\rho \frac{\Pi \parallel S}{B \vee C} \star$	$\backslash odR\{\backslash Pi\}\{[B.C]\}\{\backslash cal S\}$	
$\rho \frac{\Pi \parallel S}{B \vee C} \star$	$\backslash odR\{\backslash Pi\}\{[B.C]\}\{\backslash cal S\}$	
$\rho \frac{A}{\Delta \parallel S} \star$	$\{\{A\}\}$	$\backslash od\{\backslash odd\{\backslash odh$
	$\{\{\Delta\}\{[B.C]\}\{\backslash cal S\}\}$	$\}$
$\rho \frac{A}{\Delta \parallel S} \star$	$\{\{A\}\}$	$\backslash od\{\backslash odD\{\backslash odh$
	$\{\{\Delta\}\{[B.C]\}\{\backslash cal S\}\}$	$\}$
$\rho \frac{A}{\Delta \parallel S} \star$	$\{\{A\}\}$	$\backslash od\{\backslash oddx\{\backslash odh$
	$\{\{\Delta\}\{[B.C]\}\{\backslash cal S\}\{\the\toks0\}\}$	$\}$
$\rho \frac{A}{\Delta \parallel S} \star$	$\backslash odv\{A\}\{\Delta\}\{[B.C]\}\{\backslash cal S\}$	
$\rho \frac{A}{\Delta \parallel S} \star$	$\backslash odV\{A\}\{\Delta\}\{[B.C]\}\{\backslash cal S\}$	

$\frac{\Pi \vdash S}{A} \rho$	$B \vee C$ note	$\text{\odf{\odif{\odp}}}$ $\{\backslash Pi\}\{A\}\{\backslash cal{S}\}\}$ $\{\backslash rho\}\{[B.C]\}\{\backslash rm{note}\}$
$\frac{A}{\Delta \parallel S}$	B	$\text{\odf{\odif{\odd{\odh}}}}$ $\{A\}\}$
$\frac{B}{\rho \frac{C}{C}}$		$\{\backslash Delta\}\{B\}\{\backslash cal{S}\}$
		$\{\backslash rho\}\{C\}\{ \}$
$\frac{A}{\Delta \parallel S''}$	B	$\text{\odf{\odd{\odd{\odd{\odh}}}}}$ $\{A\}\}$
$\frac{B}{\Delta' \parallel S'}$	C	$\{\backslash Delta\}\{B\}\{\backslash cal{S}'\}$
$\frac{C}{\Delta'' \parallel S}$	D	$\{\backslash Delta'\}\{C\}\{\backslash cal{S}'\}$
$\frac{A}{\Delta \parallel S}$	D	$\{\backslash Delta''\}\{D\}\{\backslash cal{S}\}$
$\frac{A}{B}$		$\text{\odf{\odd{\odd{\odif{\odh}}}}}$ $\{A\}\}$
$\frac{\Delta \parallel C}{D}$		$\{\backslash rho\}\{B\}\{ \}$
		$\{\backslash Delta\}\{C\}\{ \}$
$\frac{A}{B}$		$\{\backslash Delta'\}\{D\}\{ \}$
$\frac{\rho \frac{A}{B}}{\rho' \frac{B}{C}}$		$\text{\odf{\odif{\odI{\odo{\odh}}}}}$ $\{A\}\}$
$\frac{\rho' \frac{A}{B}}{\rho'' \frac{B}{C}}$		$\{\backslash rho\}\{B\}\{ \}$
$\frac{\rho'' \frac{A}{B}}{\rho''' \frac{B}{D}}$		$\{\backslash rho'\}\{C\}\{ \}$
		$\{\backslash rho''\}\{D\}\{\backslash quad.\}$
$A \otimes B$		$\text{\odV{[A;B]}}$
$\frac{}{\psi_1 \parallel}$		$\{\backslash psi_1\}\{\backslash odbox{[a;b}$
$\frac{}{a \otimes b \otimes \frac{\bar{b} \otimes b}{\perp} \otimes (\bar{a} \otimes \bar{b})}$		$; \text{\odn{(-b;b)}} \text{\iU{\bot}} \{ \}$
		$; (-a;-b)$
		$] } \{ \}$
t		\odbackgroundtrue
$\frac{i \downarrow}{\frac{a}{a \wedge t} \vee \frac{\bar{a}}{t \wedge \bar{a}}}$		$\text{\odf{\odif{\odif{\odif{\odh}}}}}$ $\{\backslash ttt\}$
m		$\backslash id{[\backslash odn{ a }]$
$\frac{(a \vee t) \wedge \frac{t \vee \bar{a}}{\bar{a} \vee t}}{s}$		$=\{(\backslash ttt).a\}..\backslash odn{ -a }$
s		$=\{(\backslash ttt.-a)\}[] } \{ \}}$
$\frac{(a \vee t) \wedge \bar{a}}{\bar{a} \wedge (a \vee t)}$		$\backslash me{([a.\backslash ttt]..\backslash odn{[\backslash ttt.-a]})$
$i \uparrow$		$=\{[-a.\backslash ttt]\}[] } \{ \}}$
f		$\backslash sw{[\text{\odf{\odif{\odif{\odh}}}}$
		$\{([a.\backslash ttt].-a)\} }$
		$=\{(-a.[a.\backslash ttt])\} }$
		$\backslash sw{[\backslash odn{(-a.a)}$
		$\text{\iU{\ffff}} \{ \}..\backslash ttt]\}[] } \{ \}}$

You can suppress or enable the background colour and the frames, respectively, with the commands `\odbackgroundfalse` or `\odbackgroundtrue`, and `\odframefalse` or `\odframetrue`.

6. ATOMIC FLOWS

The following uses XY-pic (see at the beginning of this manual for instructions).

	\afid abcd
	\afidc abcd{Red}{Green}
	\afidx abcd45
	\afidxc abcd45{Red}{Green}
	\afidn
	\afId abcd
	\afIdc abcd{Red}{Green}
	\afIdx abcd45
	\afIdxc abcd45{Red}{Green}

	\afiu abcd
	\afiu abcd{Red}{Green}
	\afiu abcd32
	\afiu abcd32{Red}{Green}
	\afiuun
	\afIu abcd
	\afIuc abcd{Red}{Green}
	\afIux abcd32
	\afIux abcd32{Red}{Green}

The three macros `\aftrim`, `\aftrimabove` and `\aftrimbelow` trim extra space above and below diagrams including `\afiu` and similar macros.

	\aftrim{\af{(0,0)*{\afwun}}}
	\aftrimabove{\af{(0,0)*{\afid{}{}{}{}}}}
	\aftrimbelow{\af{(0,0)*{\afwu{}{}{}}}}

Use the macros `\afraise` and `\aflower`, typically with one atomic flow as an argument, to raise or lower a bit the atomic flow (see examples in the following).

▼ \afwd {}{}

$a \mid b$ \afwd ab

$a \mid b$ \afwdc ab{Red}

▽ \afwdn

▼ \afWd {}{}

$a \parallel b$ \afWd ab

$a \parallel b$ \afWdc ab{Red}

▼ \afwds {}{}

$a \mid b$ \afwds ab

$a \mid b$ \afwdsc ab{Red}

▼ \afwdsn

▼ \afWds {}{}

$a \mid b$ \afWds ab

$a \mid b$ \afWdsc ab{Red}

▲ \afwu {}{}

$a \mid b$ \afwu ab

$a \mid b$ \afwuc ab{Red}

△ \afwun

▲ \afWu {}{}

$a \parallel b$ \afWu ab

$a \parallel b$ \afWuc ab{Red}

▲ \afwus {}{}

$a \mid b$ \afwus ab

$a \mid b$ \afwusc ab{Red}

▲ \afwusn

▲ \afWus {}{}

$a \mid b$ \afWus ab

$a \mid b$ \afWusc ab{Red}

 \afdd {}{{}}

 a|b \afdd ab

 a|b \afddc ab{Red}

 \afDd {}{{}}

 a||b \afDd ab

 a||b \afDdc ab{Red}

 a|b \afdu {}{{}}

 a|b \afdu ab

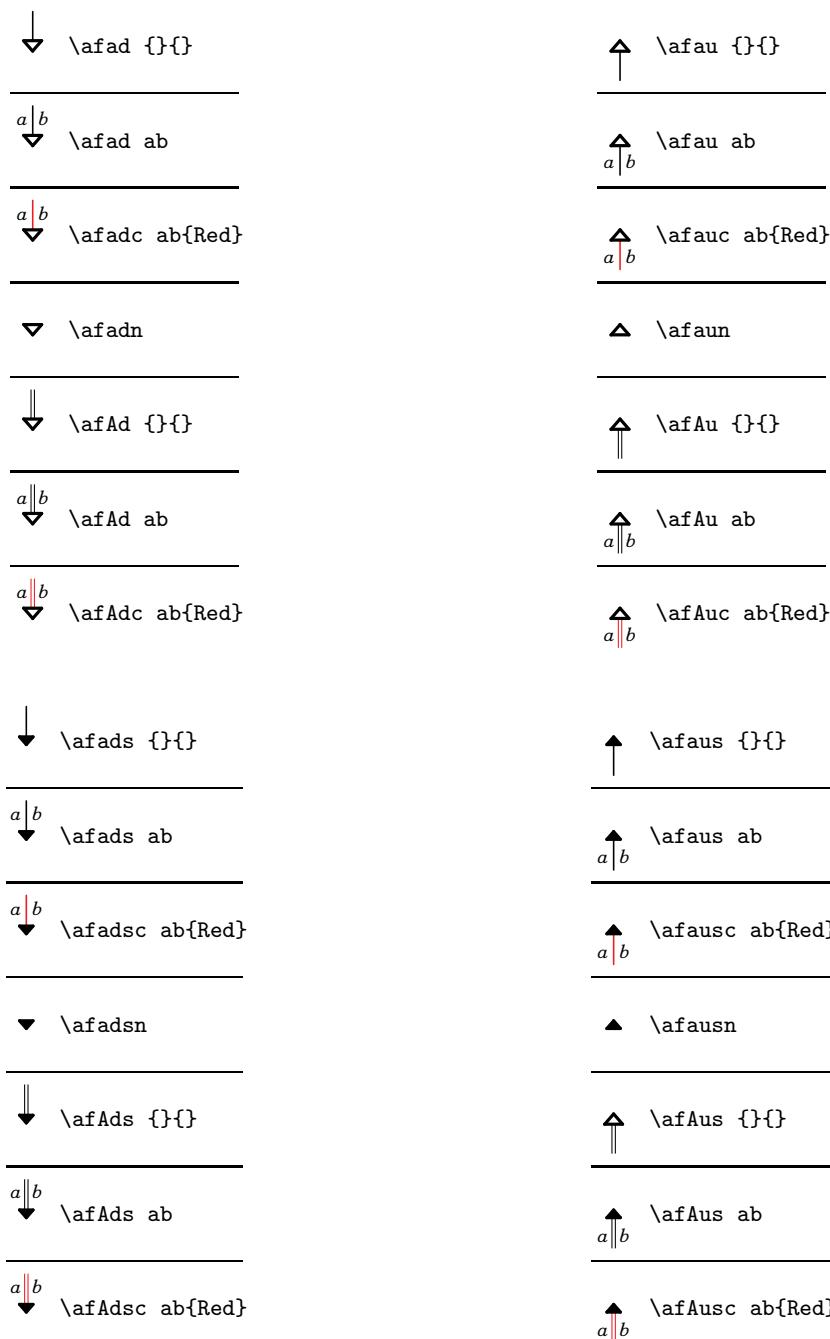
 a||b \afduc ab{Red}

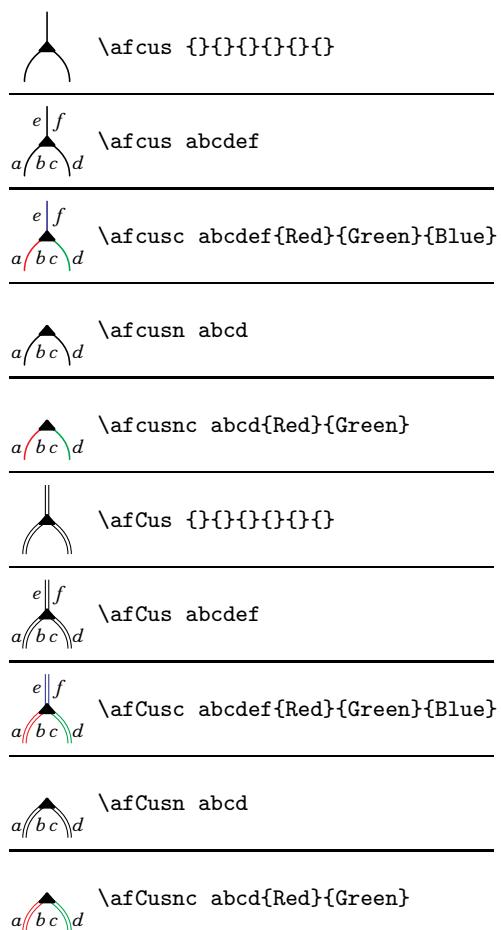
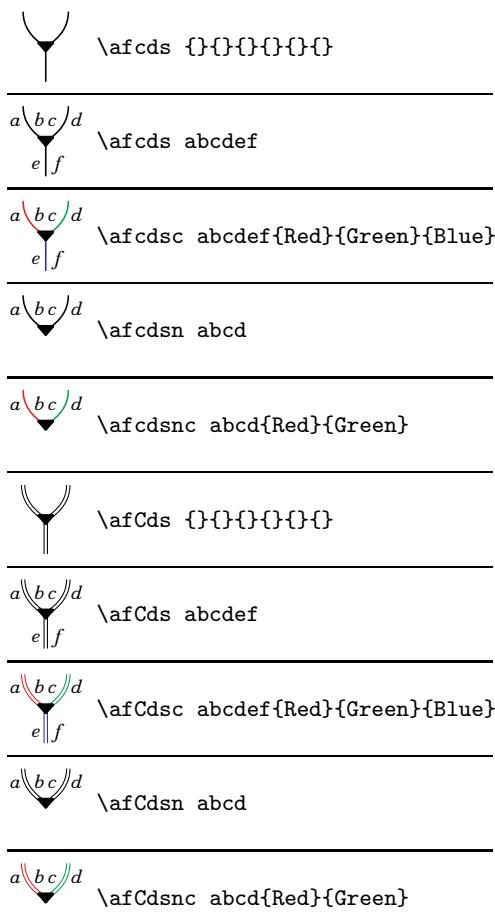
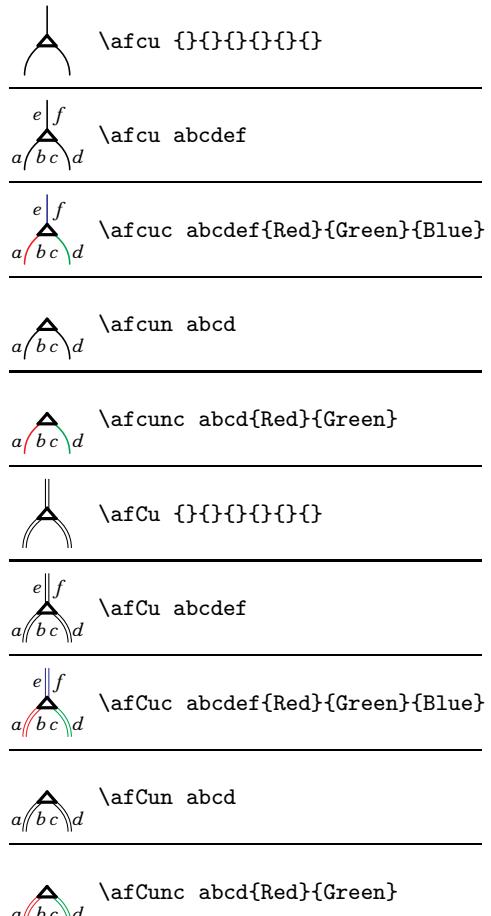
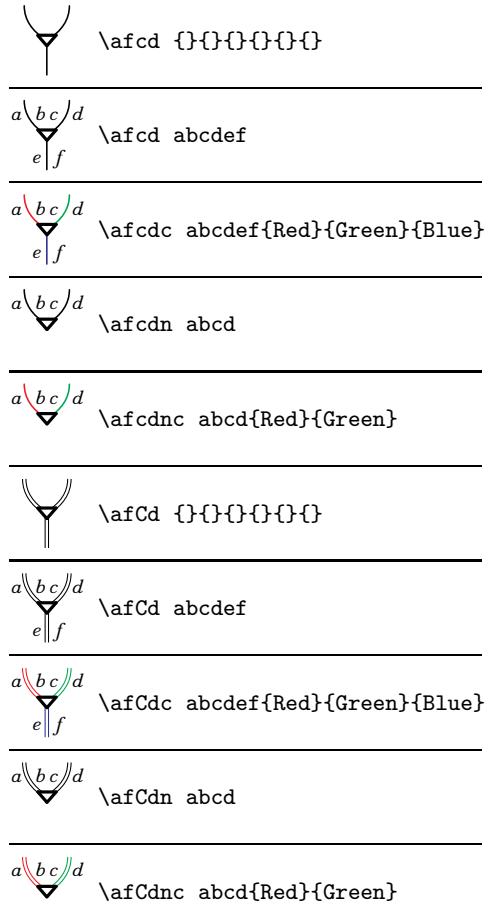
 \afDu {}{{}}

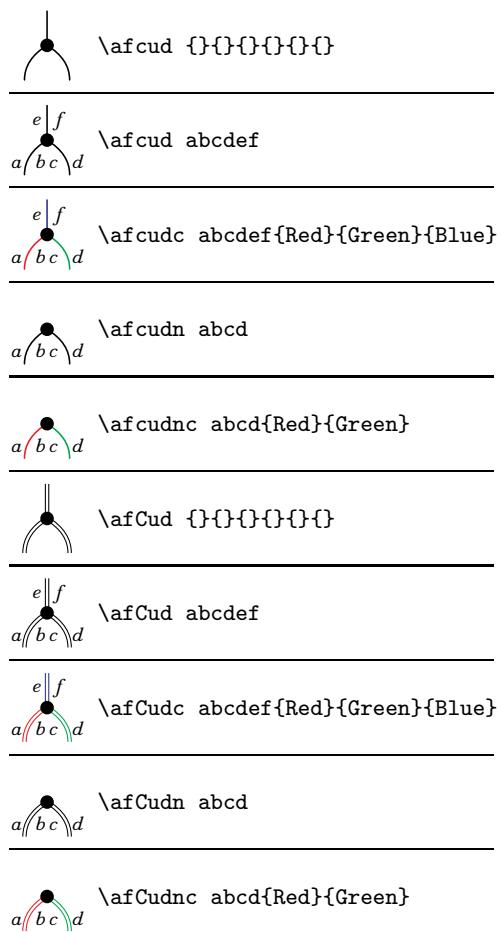
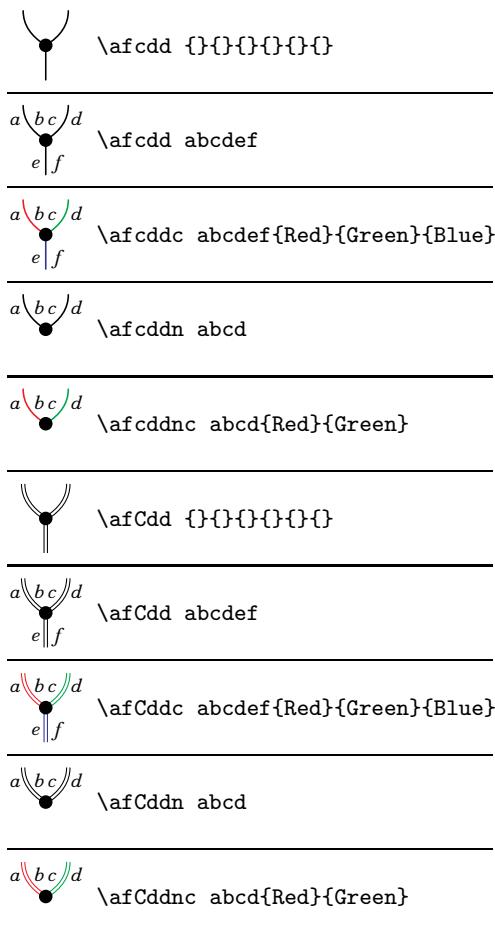
 a||b \afDu ab

 a||b \afDuc ab{Red}

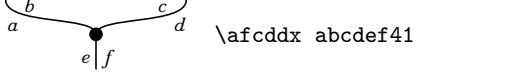
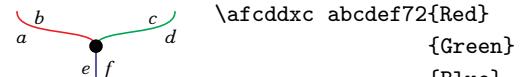
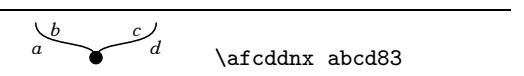
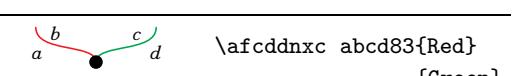
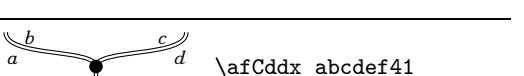
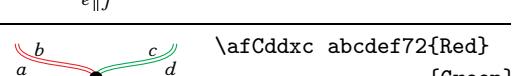
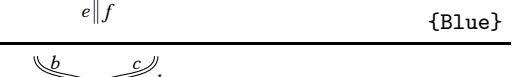
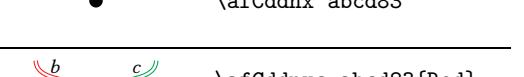
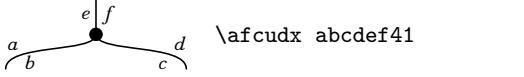
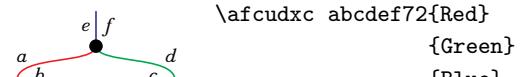
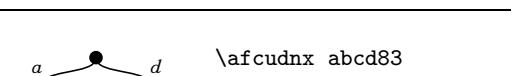
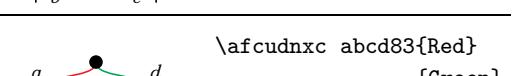
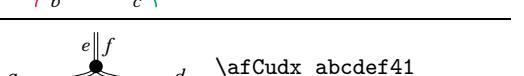
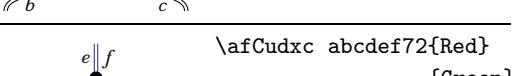
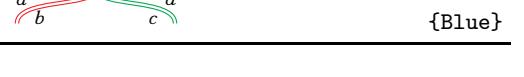
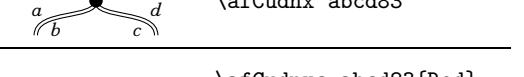
● \afd

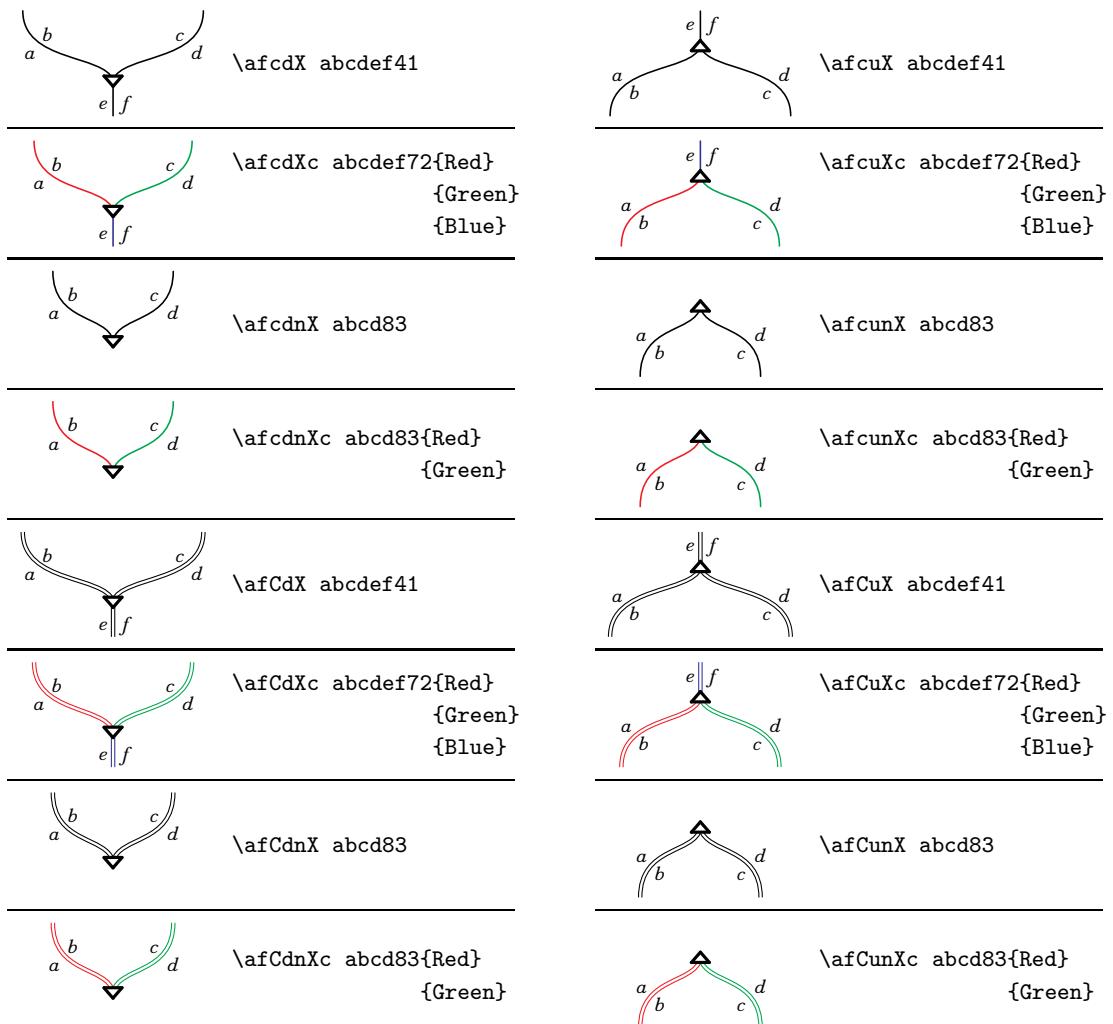


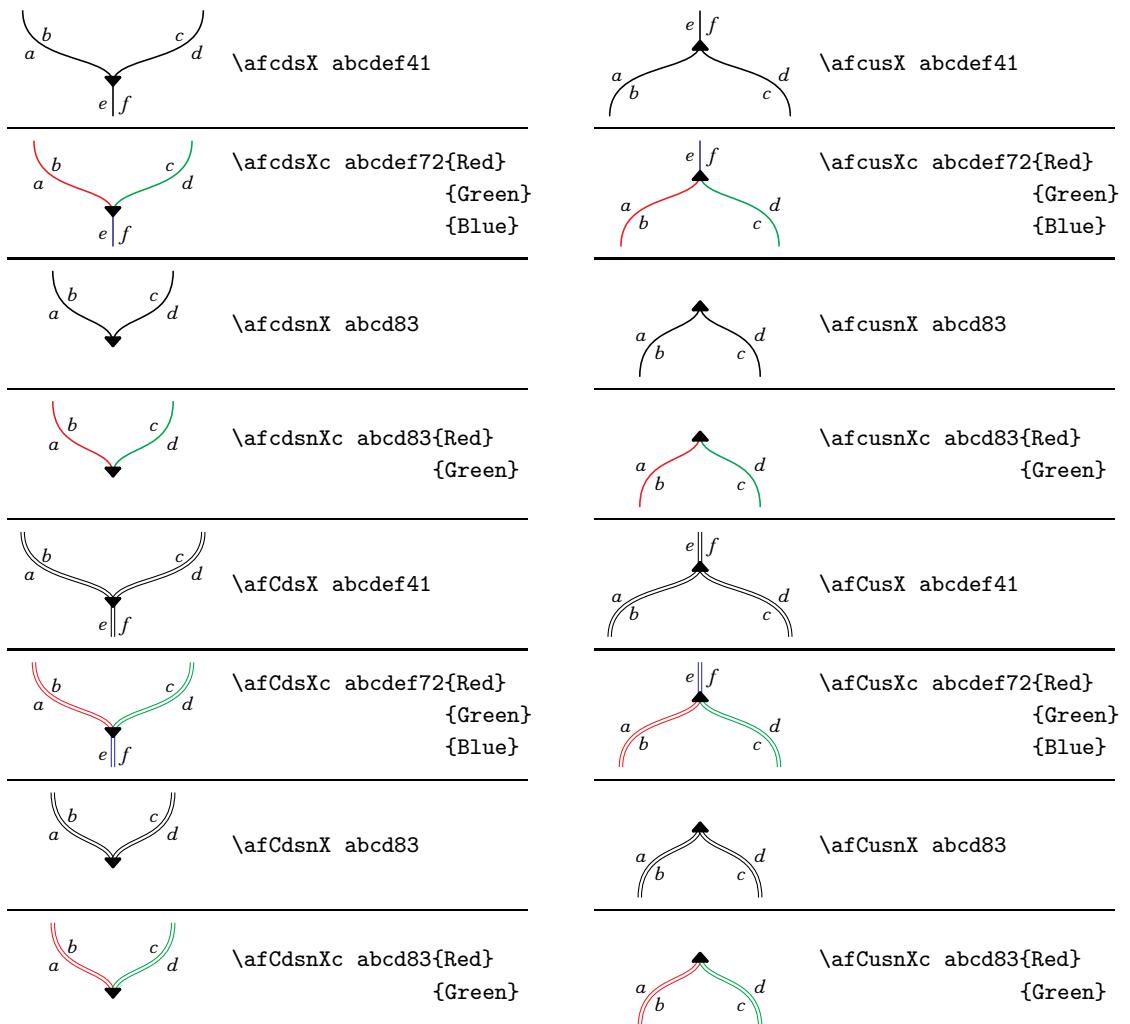


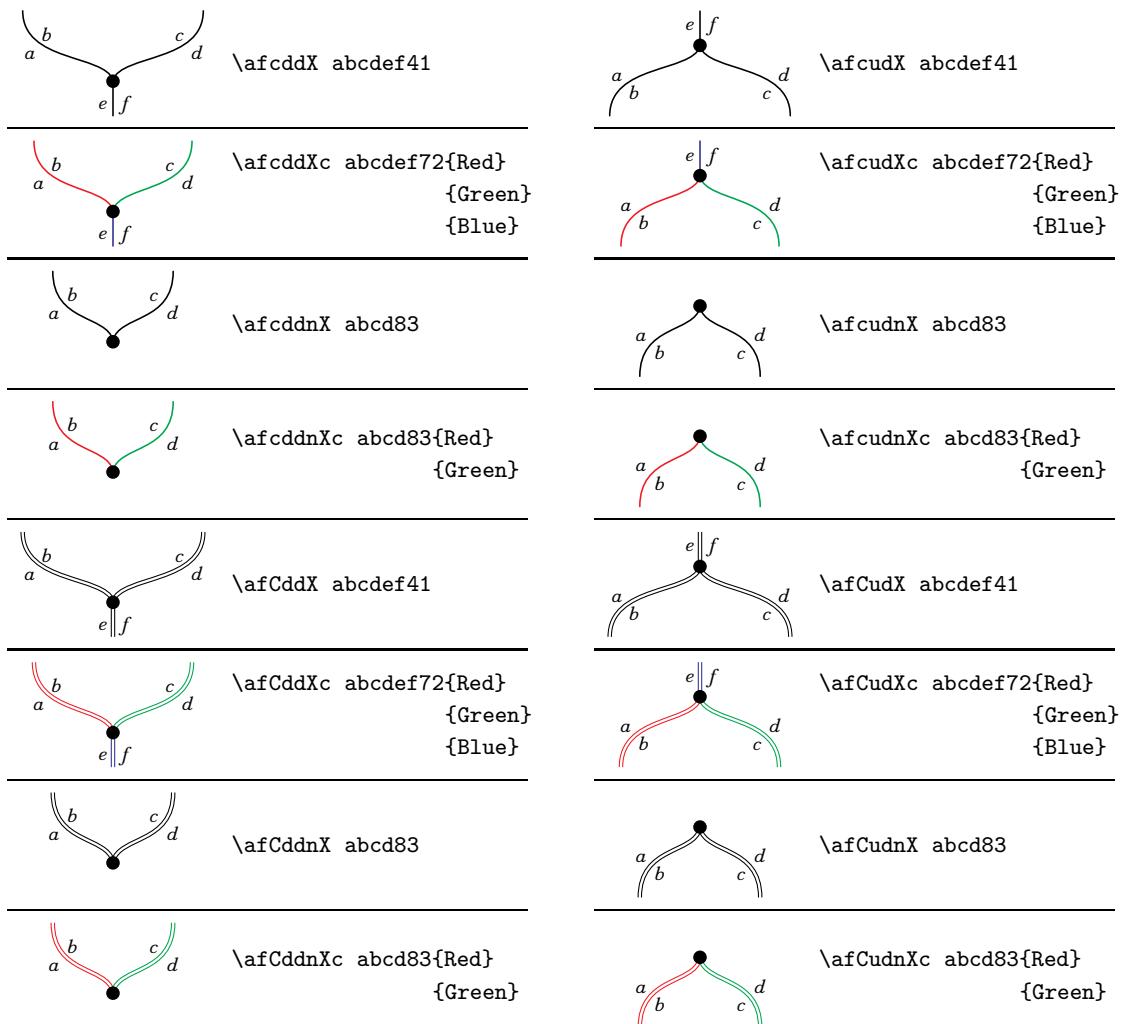


	\afcdx abcdef41
	\afcdxc abcdef72{Red} {Green} {Blue}
	\afcdnx abcd83
	\afcdnxc abcd83{Red} {Green}
	\afCdx abcdef41
	\afCdx abcdef72{Red} {Green} {Blue}
	\afCdnx abcd83
	\afCdnxc abcd83{Red} {Green}
	\afcdsx abcdef41
	\afcdsx abcdef72{Red} {Green} {Blue}
	\afcdsnx abcd83
	\afcdsnxc abcd83{Red} {Green}
	\afCdsx abcdef41
	\afCdsxc abcdef72{Red} {Green} {Blue}
	\afCdsn x abcd83
	\afCdsnxc abcd83{Red} {Green}
	\afcusx abcdef41
	\afcusxc abcdef72{Red} {Green} {Blue}
	\afcusnx abcd83
	\afcusnxc abcde83{Red} {Green}
	\afCusx abcdef41
	\afCusxc abcdef72{Red} {Green} {Blue}
	\afCusnx abcd83
	\afCusnxc abcde83{Red} {Green}

	\afcddx abcdef41
	\afcddxc abcdef72{Red} {Green} {Blue}
	\afcddnx abcd83
	\afcddnxc abcd83{Red} {Green}
	\afCddx abcdef41
	\afCddxc abcdef72{Red} {Green} {Blue}
	\afCddnx abcd83
	\afCddnxc abcd83{Red} {Green}
	\afcudx abcdef41
	\afcudxc abcdef72{Red} {Green} {Blue}
	\afcudnx abcd83
	\afcudnxc abcd83{Red} {Green}
	\afCudx abcdef41
	\afCudxc abcdef72{Red} {Green} {Blue}
	\afCudnx abcd83
	\afCudnxc abcd83{Red} {Green}

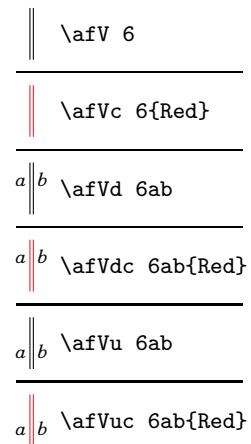
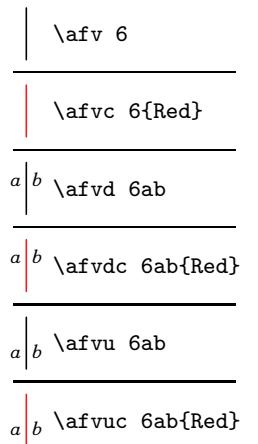




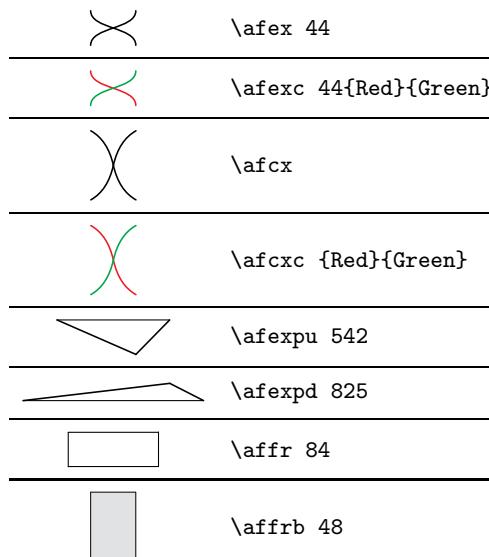
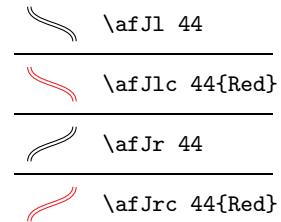
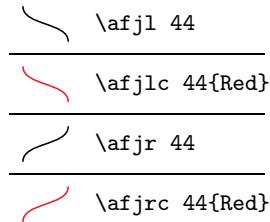


± \afbd

⊤ \afbu



: \afvdj



By default the background colour is defined as `\newxycolor{afbackground}{0 0 0 0.12 cmyk}`.

The following dimen parameters are defined:

```
\afelwidth      =425\afunit  
\afelheight    =325\afunit  
\afthickone    = 60\afunit  
\afthicktwo    = 40\afunit  
\afthickthree  =160\afunit  
\afthickfour   =120\afunit  
\afthickfive   = 40\afunit  
\aflabeldistance =220\afunit
```

The registers govern various parameters in the atomic flow elements (experiment to see which ones). You might be interested especially in changing the last one. The parameters can be globally scaled by invoking `\afsetparams` after having adjusted `\afunit`, which by default is 0.01pt.

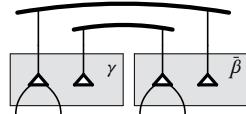
Each line starting with a coordinate pair (x, y) is an element of the figure. The coordinates are relative, so, for example, you can add or subtract the same vector to all of them and still get the same figure.

	<pre>\af{ (0,8)*{\afc{a}{a}{a}{a}}; (0,0)*{\afc{a}{a}{a}{a}}</pre>
	<pre>\af{ (2,12)*{\afc{a}{a}{a}{a}}; (8,12)*{\afc{a}{a}{a}{a}}; (0,6)*{\afv{4}}; (5,6)*{\afex{24}}; (10,6)*{\afv{4}}; (2,0)*{\afc{a}{a}{a}{a}}; (8,0)*{\afc{a}{a}{a}{a}}</pre>
	<pre>\af{ (4,4.5)*{\afvd{j}}; (0,5)*{\afc{a}{a}{a}{a}}; (2,0)*{\afiu{a}{a}{\bar{a}}}</pre>
→	<pre>\aflower{\af{ (6,6)*{\afvd{4}{}}; (3,0)*{\afiu{x}{}}32; (0,4)*{\afcd{n}{}}} \quad\quad\quad \aflower{\af{ (6,6)*{\afc{a}{a}{a}{a}}; (0,6)*{\afvd{8}{}}; (12,6)*{\afvd{8}{}}; (10,2)*{\afiu{n}}; (2,2)*{\afiu{n}}}}</pre>
→	<pre>\aflower{\af{ (6,6)*{\afVd{4}{}}; (3,0)*{\afIx{x}{}}32; (0,4)*{\afCd{n}{}}} \quad\quad\quad \aflower{\af{ (6,6)*{\afc{a}{a}{a}{a}}; (0,6)*{\afVd{8}{}}; (12,6)*{\afVd{8}{}}; (10,2)*{\afiu{n}}; (2,2)*{\afiu{n}}}}</pre>
	<pre>\af{ (4,4.5)*{\afvd{j}}; (0,5)*{\afc{a}{a}{a}{a}}; (1,2.5)*{\affr{10}{13}}; (2,0)*{\afiu{a}{a}{\bar{a}}}</pre>
	<pre>\af{ (4,16)*{\afid{}}; (0,8)*{\afc{a}{a}{a}{a}}; (5,8)*{\afjr{28}}; (2,0)*{\afiu{a}{a}{\bar{a}}}}</pre>
	<pre>{\af{ (4,16)*{\afid{}}; (0,8)*{\afc{a}{a}{a}{a}}; (6,8)*{\afv{8}}; (3,0)*{\afiu{x}{}}32}}</pre>

```

\af{
( 2 ,6 )*\{\afidc{}{}{}{}74{white}{white}\};
(-1.5,5 )*\{\afv2\};
( 5.5,5 )*\{\afv2\};
( 2 ,8 )*\{\afidx{}{}{}{}15\}4;
( 8 ,0 )*\{\affrb{10}6\};
(13 ,1.5)*{\afll{\bar\beta\beta}};
( 9.5,0 )*\{\afwu{}\}{};
( 5.5,0 )*\{\afc{}{}{}{}{}\}{};
(-3 ,0 )*\{\affrb{10}6\};
( 2 ,1.5)*{\afll\gamma\gamma};
(-1.5,0 )*\{\afwu{}\}{};
(-5.5,0 )*\{\afc{}{}{}{}{}\}{}}

```



In the next two figures, notice how vertical symmetry is achieved:

```

\af{
(10,20.5)*{\afvdj\}; (4,18)*{\afvd{12}a\}; (0,16)*{\afvd{16}a\}; (10,16)*{\afc{}{}{}{}{}\}{}{\bar a\}}; (6,10)*{\afex44\}; (12,10)*{\afv4\}; (2, 4)*{\afiu a\}{}{\bar a\}}; (10, 4)*{\afiu a\}{}{\bar a\}}}
\afinv{
(10,20.5)*{\afvdj\}; (4,18)*{\afvu{12}a\}; (0,16)*{\afvu{16}a\}; (10,16)*{\afcd{}{}{}{}{}\}{}{\bar a\}}; (6,10)*{\afex44\}; (12,10)*{\afv4\}; (2, 4)*{\afid a\}{}{\bar a\}}; (10, 4)*{\afid a\}{}{\bar a\}}}

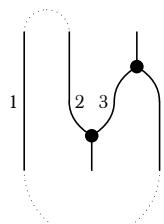
```

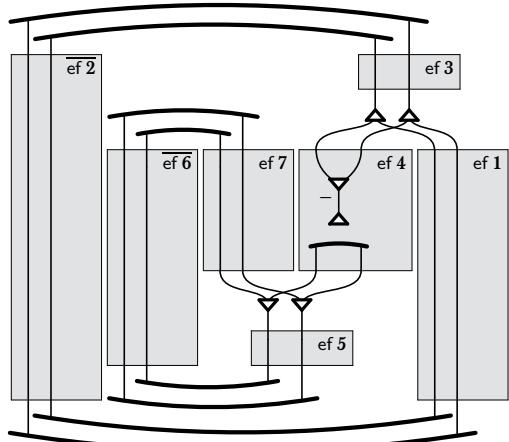
The following figure requires \xyoption{arc}:

```

\af{
(10,12)*{\afcud{}{}{}{}{}{}{}{}{}{}{}{}};
( 6, 4)*{\afcdd{}{}{}{}{}{}{}{}{}{}{}{}};
( 4,12)*{\afv8\}; (12, 4)*{\afv8\};
( 0, 8)*{\afv{16}{}{}{}{}{}{}{}{}{}{}{}{}};
( 0, 8)*{\afll11\}; ( 4, 8)*{\aflr2\};
( 8, 8)*{\afll3\}; (12, 8)*{\aflr4\};
( 0,16);(4,16),{\ellipse{.}}; (12, 0);(0, 0),{\ellipse{.}}}

```





```

\renewcommand{\one }{{{\bf \scriptstyle 1}}}
\newcommand{\two }{{{\bf \scriptstyle 2}}}
\newcommand{\three }{{{\bf \scriptstyle 3}}}
\newcommand{\four }{{{\bf \scriptstyle 4}}}
\newcommand{\five }{{{\bf \scriptstyle 5}}}
\newcommand{\six }{{{\bf \scriptstyle 6}}}
\newcommand{\seven }{{{\bf \scriptstyle 7}}}
\newcommand{\ef}{\mathop{\mathsf{ef}}\nolimits}

\af
( 1.5 , 4 )*\{ \affrb{10}{14} ; ;
( 0 , 3 )*\{ \afwu{} ; ;
( 0 , 0 )*\{ \afidn ; ;
( 6.5 , 9.5)*{ \afll{\ef\four} } ; ;
( 0 , 7 )*\{ \afcdn{}{}{}{}{} ; ;
(- 2 , - 1.5)*{ \afv3} ; ;
( 2 , - 1.5)*{ \afv3} ; ;
(- 8 , 4 )*\{ \affrb8{14} } ; ;
(-10.5 , 3 )*\{ \afv{12} } ; ;
(- 8.5 , 4 )*\{ \afv{14} } ; ;
(- 4 , 9.5)*{ \afll{\ef\seven} } ; ;
(-11.625,-16 )*\{ \afiuxc{}{}{}{}{10}4
                           {white}{white}\} ; ;
(-11.125,-18 )*\{ \afiuxc{}{}{}{}{59}{16}
                           {white}{white}\} ; ;
(- 3.25 ,-12 )*\{ \affrb94} ; ;
(- 6.25 ,-13.5)*{ \afv5} ; ;
(- 3.25 ,-14.5)*{ \afv7} ; ;
( 1.25 ,-11.5)*{ \afll{\ef\five} } ; ;
(- 6.25 ,- 7 )*\{ \afcdx{}{}{}{}{}{17}8\} ; ;
(- 3.25 ,- 7 )*\{ \afcdx{}{}{}{}{}{21}8\} ; ;
(-10.625, 26 )*\{ \afidxc{}{}{}{}{59}8
                           {white}{white}\} ; ;
(-11.125, 24 )*\{ \afidxc{}{}{}{}{101}{16}
                           {white}{white}\} ; ;
(- 8.5 ,-22 )*\{ \afiuxc{}{}{}{}{33}4
                           {white}{white}\} ; ;
(- 8.5 ,-20 )*\{ \afiuxc{}{}{}{}{59}8
                           {white}{white}\} ; ;
(-16.5 , - 1.5)*{ \affrb8{25} } ; ;
(-19 , - 3.5)*{ \afv{29} } ; ;
(-17 , - 3.5)*{ \afv{25} } ; ;
(-12.5 , 9.5)*{ \afll{\overline{\ef\six}}} ; ;
(-13.75 , 13 )*\{ \afidx{}{}{}{}{13}8\} ; ;
(-13.75 , 15 )*\{ \afidx{}{}{}{}{21}8\} ; ;
( 11 , -3.5)*{ \affrb8{29} } ; ;
( 8.5 , -4.5)*{ \afv{31} } ; ;
( 10.5 , -5.5)*{ \afv{33} } ; ;
( 15 , 9.5)*{ \afll{\ef\one} } ; ;
( 3.25 , 15 )*\{ \afcux{}{}{}{}{}{21}8\} ; ;
( 6.25 , 15 )*\{ \afcux{}{}{}{}{}{17}8\} ; ;
( 6.25 , 20 )*\{ \affrb94} ; ;
( 3.25 , 21 )*\{ \afv6} ; ;
( 6.25 , 22 )*\{ \afv8} ; ;
( 10.75 , 20.5)*{ \afll{\ef\three} } ; ;
(-25 , 2 )*\{ \affrb8{40} } ; ;
(-27.5 , 2 )*\{ \afv{48} } ; ;
(-25.5 , 2 )*\{ \afv{44} } ; ;
(-21 , 20.5)*{ \afll{\overline{\ef\two}}} ;

```

```

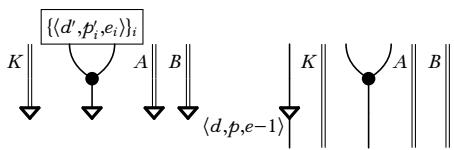
\af{
( 7,20)*{\afvc{28}{RawSienna}};
(-3, 8)*{\afvc{4}{RubineRed}};
(-2,20)="A";
"A"+( 5, 8)*{\afvc{12}{Green}};
"A"+(-3,10)*{\afcuc{}{}{}{}{}{}{}{Periwinkle}
{YellowOrange}
{OliveGreen}};
"A"+(-1, 6)*{\afcunc{}{}{}{}{}{}{}{LimeGreen}
{ProcessBlue}};
"A"+( 3, 2)*{\afiun};
"A"+(-5,-4)*{\afvc{20}{Periwinkle}};
"A"+(-3,14)*{\afwdn};
(-3,18)*{\afjlc48{LimeGreen}};
(-1,14)*{\afcunc{}{}{}{}{RubineRed}
{Aquamarine}};
( 1,10)*{\afcunc{}{}{}{}{Lavender}
{SpringGreen}};
( 5, 6)*{\afiun};
(-1, 6)*{\afwun};
( 5,38)*{\afcuc{}{}{}{}{}{}{}{Green}
{RawSienna}
{Red}};
(-5, 2)*{\afcdc{}{}{}{}{}{}{}{Periwinkle}
{RubineRed}
{DarkOrchid}}}

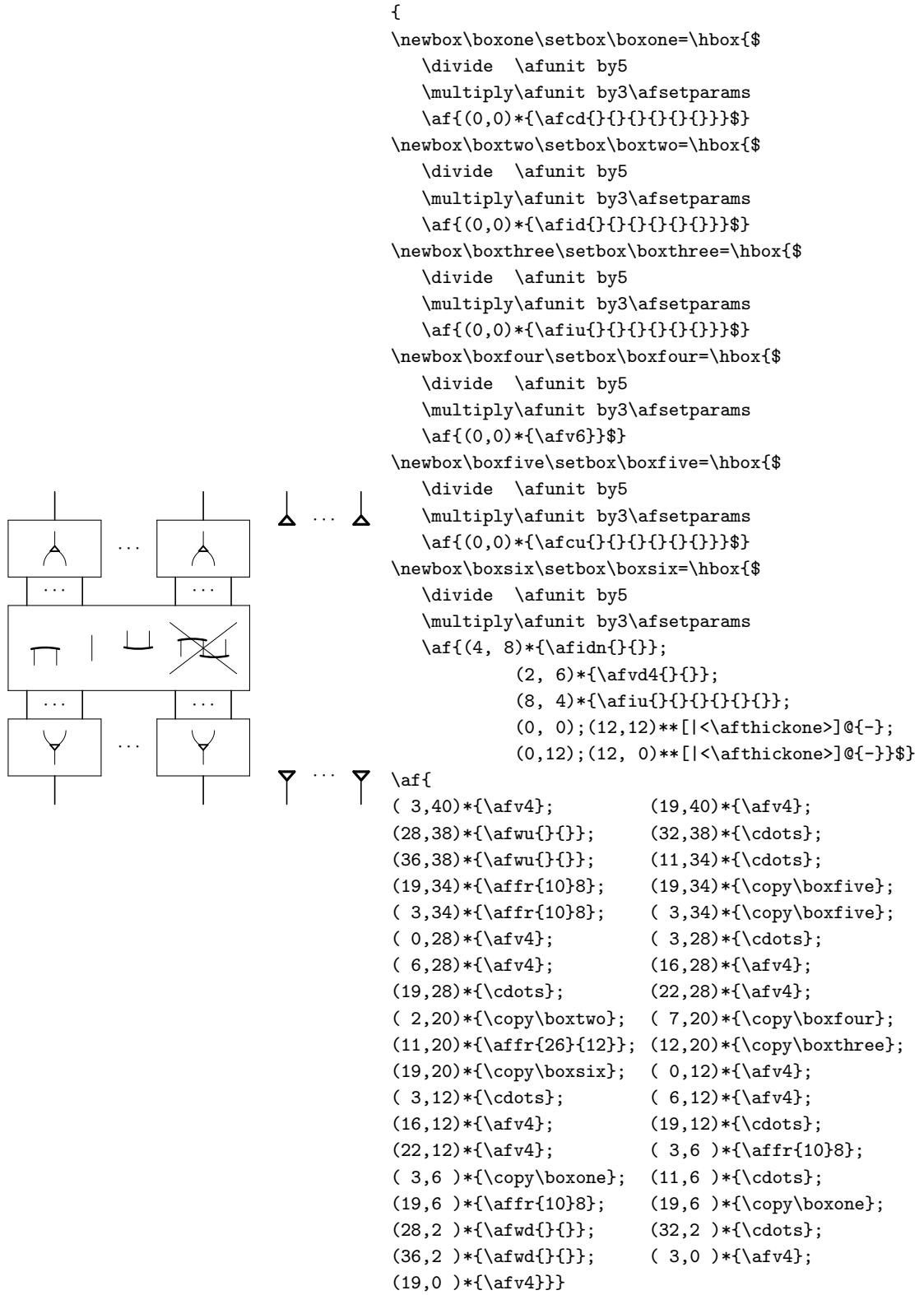
```

```

\af{
(-1 , 4 )*\{\afAd{}{}\} ;
(-1 ,10 )*\{\afVd4K{}\} ;
( 4.5,14 )*\{\affr94\} ;
( 4.5,14 )*\{\strut\scriptstyle
\{{\langle d',p'_i,e_i \rangle}_i\}
( 4.5, 8 )*\{\afcdd{}{}{}{}{}{}\} ;
( 4.5, 4 )*\{\afadn\} ;
(10 , 4 )*\{\afAd{}{}\} ;
(10 ,10 )*\{\afVd4A{}\} ;
(13 , 4 )*\{\afAd{}{}\} ;
(13 ,10 )*\{\afVd4B{}\} ;
(22 , 2 )*\{\afv4\} ;
(22 , 4 )*\{\afad{}\} ;
(22 ,10 )*\{\afv4\} ;
(22 , 2.5)*{\afll{\strut\scriptstyle
\langle d,p,e-1 \rangle}} ;
(25 , 6 )*\{\afVd{12}K{}\} ;
(29 , 2 )*\{\afv4\} ;
(29 , 8 )*\{\afcdd{}{}{}{}{}{}\} ;
(33 , 6 )*\{\afVd{12}A{}\} ;
(36 , 6 )*\{\afVd{12}B{}\} ;
}

```





7. SMASH MACROS

Compare:

α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$	α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$	α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$ α $\rho \parallel^S \beta$
		aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$
		aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vldownsmash$ ${\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vlupsmash{\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$
		aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa aaa $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$ $\$\\vlder{\\rho}{\\cal S}{\\beta}{\\alpha}$

8. ACKNOWLEDGEMENTS

Many thanks to Ross Moore for suggestions leading to `\vlupdate` and to Peter Wilson for allowing me to include the code of `ifmtarg.sty`.

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